



INTERNATIONAL STANDARD ISO/IEC 9945-2:2003
TECHNICAL CORRIGENDUM 1

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Information technology — Portable Operating System Interface (POSIX®) —

Part 2: System Interfaces

TECHNICAL CORRIGENDUM 1

Technologies de l'information — Interface pour la portabilité des systèmes (POSIX®) —

Partie 2: Interfaces systèmes

RECTIFICATIF TECHNIQUE 1

Technical Corrigendum 1 to ISO/IEC 9945-2:2003 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 22, *Programming languages, their environments and system software interfaces*.

1 Scope

This technical corrigendum addresses issues raised in defect reports and interpretation requests submitted up to 14 August 2003, and that meet all of the following criteria:

- a. They are in the scope of the approved International Standard.
- b. They contain no new APIs (functions/utilities), however, they may add enumeration symbols, non-function # defines, and reserve additional namespaces.
- c. They address contradictions between different parts of the International Standard, or add consistency between it and overriding International Standards, or address security-related problems.

2 Changes to ISO/IEC 9945-2

Change Number: XSH/TC1/1 [XBD ERN 20]

On Page: xxx Line: "ISO/IEC 8859" Section: Referenced Documents

Add after line starting "Part 10":

"Part 11: Latin/Thai Alphabet"

Add after line starting "Part 15":

"Part 16: Latin Alphabet No. 10"

Change Number: XSH/TC1/2 [XSH ERN 167, TC2d5 ERN 9]

On Page: 15, 19 Line: 580-583, 713-714 Section: 2.2.2

On lines 580-583

Change From:

"Implementations may add symbols to the headers shown in the following table, provided the identifiers for those symbols begin with the corresponding reserved prefixes in the following table, and do not use the reserved prefixes `posix_`, `POSIX_`, or `_POSIX_`."

To:

"Implementations may add symbols to the headers shown in the following table, provided the identifiers for those symbols either:

1. begin with one of the corresponding reserved prefixes in the table; or
2. have one of the corresponding complete names in the table; or
3. end in the string indicated as a reserved suffix in the table and do not use the reserved prefixes `posix_`, `POSIX_`, or `_POSIX_`, as long as the reserved suffix is in that part of the name considered significant by the implementation.

Symbols that use the reserved prefix `_POSIX_` may be made visible by implementations in any header defined by IEEE Std 1003.1-2001."

On lines 713-714

Change From:

"1. All identifiers that begin with an underscore and either an uppercase letter or another underscore are always reserved for any use by the implementation."

To:

"1. With the exception of identifiers beginning with the prefix `_POSIX_`, all identifiers that begin with an underscore and either an uppercase letter or another underscore are always reserved for any use by the implementation."

Rationale:

This change permits implementations to have symbols with the prefix `_POSIX_` visible in any header. A change in the 2003 edition had disallowed this.

Change Number: XSH/TC1/3 [XSH ERN 1]

On Page: 18 Line: 668 Section: 2.2.2

Add the following line into the table after line 668

"<math.h> FP_[A-Z]"

Rationale:

This text is added for consistency with the `math.h` page in the Base Definitions volume which states "Additional implementation-defined floating-point classifications, with macro definitions beginning with `FP_` and an uppercase letter, may also be specified by the implementation."

Change Number: XSH/TC1/4 [XSH ERN 99]

On Page: 33 Line: 1337 Section: 2.4.3

Add "socketmark()" to the table in 2.4.3 (the list of functions that shall be either reentrant or non-interruptible by signals and shall be async-signal-safe).

Rationale:

This was an omission when the socketmark() function was added.

Change Number: XSH/TC1/5 [XSH ERN 109]

On Page: 55 Line: 2267 Section: 2.9.5.2

Add "fdatasync()" into the table of functions after "fcntl()".

Rationale:

This was an omission. Since fdatsync() is often implemented in terms of fsync() it should also be in this list.

Change Number: XSH/TC1/6 [XSH ERN 76]

On Page: 56 Line: 2277 Section: 2.9.5.2 Cancellation Points

Insert into the list of functions that "may have" a cancellation point (the second list of functions in this section) in the appropriate order:

"access
 asctime
 asctime_r
 ctime
 ctime_r
 ffmtmsg
 fpathconf
 fstat
 getaddrinfo
 gethostid
 getnameinfo
 getopt [with footnote if opterr is nonzero]
 link
 localtime
 localtime_r
 lstat
 mktime
 pathconf
 posix_openpt
 stat
 strerror_r
 strptime
 symlink
 sync
 tzset
 wcsftime
 wordexp"

Change Number: XSH/TC1/7 [XSH ERN 77]

On Page: 57 Line: 2336-2353 Section: 2.9.5.3

Change From:

“Each thread maintains a list of cancellation cleanup handlers. The programmer uses the `pthread_cleanup_push()` and `pthread_cleanup_pop()` functions to place routines on and remove routines from this list. When a cancellation request is acted upon, the routines in the list are invoked one by one in LIFO sequence; that is, the last routine pushed onto the list (Last In) is the first to be invoked (First Out). The thread invokes the cancellation cleanup handler with cancellation disabled until the last cancellation cleanup handler returns. When the cancellation cleanup handler for a scope is invoked, the storage for that scope remains valid. If the last cancellation cleanup handler returns, thread execution is terminated and a status of `PTHREAD_CANCELED` is made available to any threads joining with the target. The symbolic constant `PTHREAD_CANCELED` expands to a constant expression of type `(void *)` whose value matches no pointer to an object in memory nor the value `NULL`. The cancellation cleanup handlers are also invoked when the thread calls `pthread_exit()`.

A side effect of acting upon a cancellation request while in a condition variable wait is that the mutex is re-acquired before calling the first cancellation cleanup handler. In addition, the thread is no longer considered to be waiting for the condition and the thread shall not have consumed any pending condition signals on the condition.

A cancellation cleanup handler cannot exit via `longjmp()` or `siglongjmp()`.”

To:

Each thread maintains a list of cancellation cleanup handlers. The programmer uses the `pthread_cleanup_push()` and `pthread_cleanup_pop()` functions to place routines on and remove routines from this list. When a cancellation request is acted upon, or when a thread calls `pthread_exit()`, the thread first disables cancellation by setting its cancelability state to `PTHREAD_CANCEL_DISABLE` and its cancelability type to `PTHREAD_CANCEL_DEFERRED`. The cancelability state shall remain set to `PTHREAD_CANCEL_DISABLE` until the thread has terminated. The behavior is undefined if a cancellation cleanup handler or thread-specific data destructor routine changes the cancelability state to `PTHREAD_CANCEL_ENABLE`. The routines in the thread's list of cancellation cleanup handlers are invoked one by one in LIFO sequence; that is, the last routine pushed onto the list (Last In) is the first to be invoked (First Out). When the cancellation cleanup handler for a scope is invoked, the storage for that scope remains valid. If the last cancellation cleanup handler returns, thread-specific data destructors (if any) associated with thread-specific data keys for which the thread has non-`NULL` values will be run, in unspecified order, as described for `pthread_key_create()`.

After all cancellation cleanup handlers and thread-specific data destructors have returned, thread execution is terminated. If the thread has terminated because of a call to `pthread_exit()`, the `value_ptr` argument is made available to any threads joining with the target. If the thread has terminated by acting on a cancellation request, a status of `PTHREAD_CANCELED` is made available to any threads joining with the target. The symbolic constant `PTHREAD_CANCELED` expands to a constant expression of type `(void *)` whose value matches no pointer to an object in memory nor the value `NULL`.

A side effect of acting upon a cancellation request while in a condition variable wait is that the mutex is re-acquired before calling the first cancellation cleanup handler. In addition, the thread is no longer considered to be waiting for the condition and the thread shall not have consumed any pending condition signals on the condition.

A cancellation cleanup handler cannot exit via `longjmp()` or `siglongjmp()`.”

Change Number: XSH/TC1/8 [XSH ERN 89]

On Page: 58 Line: 2379 Section: Threads

Insert new section:

“2.9.8 Use of Application-Managed Thread Stacks

An “application-managed thread stack” is a region of memory allocated by the application, for example, memory returned by the `malloc()` or `mmap()` functions, and designated as a stack through the act of passing an address related to that memory as the “`stackaddr`” argument to `pthread_attr_setstackaddr()` (obsolete) or by passing the address and size of the stack, respectively, as the `stackaddr` and `stacksize` arguments to `pthread_attr_setstack()`. Application-managed stacks allow the application to precisely control the placement and size of a stack. The application grants to the implementation permanent ownership of and control over the application-managed stack when the attributes object in which the stack or `stackaddr` attribute has been set is used, either by presenting that attributes object as the ‘`attr`’ argument in a call to `pthread_create()` that completes successfully, or by storing a pointer to the attributes object in the `sigev_notify_attributes` member of a ‘`struct sigevent`’ and passing that ‘`struct sigevent`’ to a function accepting such argument that completes successfully. The application may thereafter utilize the memory within the stack only within the normal context

of stack usage within or properly synchronized with a thread that has been scheduled by the implementation with stack pointer value(s) that are within the range of that stack. In particular, the region of memory cannot be freed, nor can it be later specified as the stack for another thread.

When specifying an attributes object with an application-managed stack through the `sigev_notify_attributes` member of a 'struct sigevent', the results are undefined if the requested signal is generated multiple times (as for a repeating timer).

Until an attributes object in which the `stack` or `stackaddr` attribute has been set is used, the application retains ownership of and control over the memory allocated to the stack. It may free or reuse the memory as long as it either deletes the attributes object, or before using the attributes object replaces the stack by making an additional call to the same function, either `pthread_attr_setstackaddr()` or `pthread_attr_setstack()`, that was used originally to designate the stack. There is no mechanism to retract the reference to an application-managed stack by an existing attributes object.

Once an attributes object with an application-managed stack has been used, that attributes object cannot be used again by a subsequent call to `pthread_create()` or any function accepting a 'struct sigevent' with `sigev_notify_attributes` containing a pointer to the attributes object, without designating an unused application-managed stack by making an additional call to the function originally used to define the stack, `pthread_attr_setstack()` or `pthread_attr_setstackaddr()`."

Change Number: XSH/TC1/9 [XSH ERN 138]

On Page: 92 Line: 3558-3559 Section: abort

In the APPLICATION USAGE section

Change From:

"Catching the signal is intended to provide the application writer with a portable means to abort processing, free from possible interference from any implementation-defined functions."

To:

"Catching the signal is intended to provide the application writer with a portable means to abort processing, free from possible interference from any implementation-supplied functions."

Change Number: XSH/TC1/10 [XSH ERN 2]

On Page: 105 Line: 3919 Section: `aio_cancel`

In the RETURN VALUE section

Change From:

"The `aio_cancel()` function shall return the value `AIO_CANCELED` to the calling process if the requested operation(s) were canceled."

To:

"The `aio_cancel()` function shall return the value `AIO_CANCELED` if the requested operation(s) were canceled."

Rationale:

The use of the term "to the calling process" was unnecessary and incorrect.

Change Number: XSH/TC1/11 [XSH ERN 3]

On Page: 108 Line: 4022 Section: `aio_fsync`

In the RETURN VALUE section

Change From:

"The `aio_fsync()` function shall return the value 0 to the calling process if the I/O operation is successfully queued;"

To:

"The `aio_fsync()` function shall return the value 0 if the I/O operation is successfully queued;"

Rationale:

The use of the term "to the calling process" was unnecessary and incorrect.

Change Number: XSH/TC1/12 [XSH ERN 4,5]

On Page: 110 Line: 4065,4088 Section: `aio_read`

In the DESCRIPTION section

Change From:

"If prioritized I/O is supported for this file, then the asynchronous operation shall be submitted at a priority equal to the scheduling priority of the process minus `aio_cbp->aio_reqprio`."

To:
"[PIO] If prioritized I/O is supported for this file, then the asynchronous operation shall be submitted at a priority equal to a base scheduling priority minus aiocbp->reqprio.

If Thread Priority Scheduling is not supported then the base scheduling priority is that of the calling process, [PIO TPS] otherwise the base scheduling priority is that of the calling thread. [/end shading]"

Rationale:

The previous wording did not take threads into account.

In the RETURN VALUE section

Change From:

"The aio_read() function shall return the value zero to the calling process if the I/O operation is successfully queued;"

To:

"The aio_read() function shall return the value zero if the I/O operation is successfully queued;"

Rationale:

The use of the term "to the calling process" was unnecessary and incorrect.

Change Number: XSH/TC1/13 [XSH ERN 83]

On Page: 111 Line: 4101-4103 Section: aio_read

In the ERRORS section

Change From:

"[EINVAL] The file offset value implied by aiocbp->aio_offset would be invalid, aiocbp->aio_reqprio is not a valid value, or aiocbp->aio_nbytes is an invalid value."

To:

"[EINVAL] The file offset value implied by aiocbp->aio_offset would be invalid, [PIO]aiocbp->aio_reqprio is not a valid value[/PIO], or aiocbp->aio_nbytes is an invalid value."

Rationale:

Detection of an [EINVAL] error for a invalid value of aiocbp->aio_reqprio should only be required if _POSIX_PRIORITIZED_IO is supported.

Change Number: XSH/TC1/14 [XSH ERN 6,7]

On Page: 116 Line: 4252,4276 Section: aio_write

In the DESCRIPTION section

Change From:

"If prioritized I/O is supported for this file, then the asynchronous operation shall be submitted at a priority equal to the scheduling priority of the process minus aiocbp->aio_reqprio."

To:

"[PIO] If prioritized I/O is supported for this file, then the asynchronous operation shall be submitted at a priority equal to a base scheduling priority minus aiocbp->reqprio. If Thread Priority Scheduling is not supported then the base scheduling priority is that of the calling process, [PIO TPS] otherwise the base scheduling priority is that of the calling thread. [/end shading]"

Rationale:

The previous wording did not take threads into account.

In the RETURN VALUE section

Change From:

"The aio_write () function shall return the value zero to the calling process if the I/O operation is successfully queued;"

To:

"The aio_write () function shall return the value zero if the I/O operation is successfully queued;"

Rationale:

The use of the term "to the calling process" was unnecessary and incorrect.

Change Number: XSH/TC1/15 [XSH ERN 83]

On Page: 117 Line: 4290-4292 Section: aio_write

In the ERRORS section

Change From:

"[EINVAL] The file offset value implied by aiocbp->aio_offset would be invalid, aiocbp->aio_reqprio is not a valid value, or aiocbp->aio_nbytes is an invalid value."

To:

"[EINVAL] The file offset value implied by aiocbp->aiio_offset would be invalid, [PIO]aiocbp->aiio_reqprio is not a valid value[/PIO], or aiocbp->aiio_nbytes is an invalid value."

Rationale:

Detection of an [EINVAL] error for a invalid value of aiocbp->aiio_reqprio should only be required if _POSIX_PRIORITIZED_IO is supported.

Change Number: XSH/TC1/16 [XSH ERN 139]

On Page: 119-120 Line: 4375-4376 Section: alarm

In the RATIONALE section

Change From:

"In some implementations, including 4.3 BSD, very large values of the seconds argument are silently rounded down to an implementation-defined maximum value"

To:

"In some implementations, including 4.3 BSD, very large values of the seconds argument are silently rounded down to an implementation-specific maximum value"

Change Number: XSH/TC1/17 [XSH ERN 115]

On Page: 122 Line: 4440 Section: asctime

In the RETURN VALUE section

Change From:

"Upon successful completion, asctime() shall return a pointer to the string."

To:

"Upon successful completion, asctime() shall return a pointer to the string. [CX] If the function is unsuccessful, it shall return NULL. [/CX]"

Change Number: XSH/TC1/18 [XSH ERN 123]

On Page: 132 Line: 4738 Section: atan2

In the EXAMPLES section

Change From:

"None."

To:

"Converting Cartesian to Polar coordinates system

The function below uses atan2() to convert a 2d vector expressed in cartesian coordinates (x,y) to the polar coordinates (rho,theta). There are other ways to compute the angle theta, using asin(), acos() or atan(). However, atan2() presents here two advantages:

- the angle's quadrant is automatically determined.
- the singular cases (0,y) are taken into account.

Finally, this example uses hypot() rather than sqrt() since it is better for special cases, see hypot() for more information.

```
#include <math.h>
```

```
void
```

```
cartesian_to_polar(const double x, const double y,
```

```
                  double *rho, double *theta
```

```
)
```

```
{
```

```
    *rho    = hypot (x,y); /* better than sqrt(x*x+y*y) */
```

```
    *theta  = atan2 (y,x);
```

```
} "
```

Change Number: XSH/TC1/19 [XSH ERN 8]

On Page: 137 Line: 4876 Section: atexit

Add to end of the APPLICATION USAGE section:

"Since the behavior is undefined if the exit() function is called more than once, portable applications calling atexit() must ensure that the exit() function is not called at normal process termination when all functions registered by the atexit() function are called. All functions registered by the atexit() function are called at normal process termination, which occurs by a call to the exit() function or a return from main() or on the last thread termination, when the behavior is as if the implementation called exit() with a zero argument at thread termination time. If, at normal process termination, a function registered by the atexit() function is called and a portable application needs to stop further exit() processing, it must call the _exit() function or the _Exit() function or one of the functions which cause abnormal process termination."

Rationale:

Additional clarifications for application usage are added.

Change Number: XSH/TC1/20 [XSH ERN 140]

On Page: 142 Line: 5007-5009 Section: basename

In the DESCRIPTION section

Change From:

"If the string consists entirely of the '/' character, basename() shall return a pointer to the string "/". If the string is exactly "/", it is implementation-defined whether '/' or "/" is returned."

To:

"If the string pointed to by path consists entirely of the '/' character, basename() shall return a pointer to the string "/". If the string pointed to by path is exactly "/", it is implementation-defined whether "/" or "/" is returned."

Change Number: XSH/TC1/21 [XSH ERN 138]

On Page: 259 Line: 8533-8534 Section: dlopen

In the DESCRIPTION section

Change From:

"If neither RTLD_GLOBAL nor RTLD_LOCAL are specified, then an implementation-defined default behavior shall be applied."

To:

"If neither RTLD_GLOBAL nor RTLD_LOCAL are specified, then the default behavior is unspecified."

Change Number: XSH/TC1/22 [XSH ERN 120]

On Page: 289 Line: 9418 Section: erf

In the EXAMPLES section

Change From:

"None."

To:

"Computing the probability for a normal variate

This example shows how to use erf() to compute the probability that a normal variate assumes a value in the range [x1,x2] with x1<=x2. This example uses the constant M_SQRT1_2 which is an XSI extension.

```
#include <math.h>
```

```
double
```

```
Phi(const double x1, const double x2)
```

```
{
```

```
    return ( erf(x2*M_SQRT1_2) - erf(x1*M_SQRT1_2) ) / 2;
```

```
} "
```


Change Number: XSH/TC1/23 [XSH ERN 132]

On Page: 294 Line: 9526 Section: errno

In the DESCRIPTION section

Add at the end of the second paragraph:

"The setting of errno after a successful call to a function is unspecified unless the description of that function specifies that errno shall not be modified".

Change Number: XSH/TC1/24 [XSH ERN 45,46,63]

On Page: 297-299 Line: 9637 Section: exec

In the DESCRIPTION section

On line 9637

Change From:

"The state of the floating-point environment in the new process image shall be set to the default."

To:

"The state of the floating-point environment in the new process image [THR]or in the initial thread of the new process image[THR] shall be set to the default."

On line 9650

Change From:

"After a successful call to any of the exec functions, any functions previously registered by atexit() are no longer registered."

To:

"After a successful call to any of the exec functions, any functions previously registered by atexit() [THR]or pthread_atfork()[THR] are no longer registered."

On lines 9675-9677

Change From:

"[PS]For the SCHED_FIFO and SCHED_RR scheduling policies, the policy and priority settings shall not be changed by a call to an exec function. For other scheduling policies, the policy and priority settings on exec are implementation-defined.[/PS]"

To:

"When the calling process image does not use the SCHED_FIFO, SCHED_RR, or SCHED_SPORADIC scheduling policies, the scheduling policy and parameters of the new process image and the initial thread in that new process image are implementation defined. [PS]When the calling process image uses the SCHED_FIFO, SCHED_RR, or SCHED_SPORADIC scheduling policies, the process policy and scheduling parameter settings shall not be changed by a call to an exec function.[/PS] [TPS]The initial thread in the new process image shall inherit the process scheduling policy and parameters. It shall have the default system contention scope, but shall inherit its allocation domain from the calling process image.[/TPS]"

Insert the following after line 9701

(before the list that begins "The new process shall inherit at least..."):

"[THR]The thread id of the initial thread in the new process image is unspecified. The size and location of the stack on which the initial thread in the new process image runs is unspecified. The initial thread in the new process image shall have its cancellation type set to PTHREAD_CANCEL_DEFERRED and its cancellation state set to PTHREAD_CANCEL_ENABLED. The initial thread in the new process image shall have all thread-specific data values set to NULL and all thread-specific data keys shall be removed by the call to exec without running destructors. The initial thread in the new process image shall be joinable, as if created with the detachstate attribute set to PTHREAD_CREATE_JOINABLE.[THR]"

Add after line 9722

"The initial thread of the new process shall inherit at least the following attributes from the calling thread:

- * Signal mask (see sigprocmask() and pthread_sigmask())

- * Pending signals (see sigpending())"

On lines 9723-9725

Change From:

"All other process attributes defined in this volume of IEEE Std 1003.1-2001 shall be the same in the new and old process images. The inheritance of process attributes not defined by this volume of IEEE Std 1003.1-2001 is implementation-defined."

To:

"All other process attributes defined in this volume of IEEE Std 1003.1-2001 shall be inherited in the new process image from the old process image. All other thread attributes defined in this volume of IEEE Std 1003.1-2001 shall be inherited in the initial thread in the new process image from the calling thread in the old process image. The inheritance of process or thread attributes not defined by this volume of IEEE Std 1003.1-2001 is implementation-defined."

On lines 9726-9728

Change From:

"A call to any exec function from a process with more than one thread shall result in all threads being terminated and the new executable image being loaded and executed. No destructor functions shall be called."

To:

"[THR]A call to any exec function from a process with more than one thread shall result in all threads being terminated and the new executable image being loaded and executed. No destructor functions or cleanup handlers shall be called.[/THR]"

Change Number: XSH/TC1/25 [XSH ERN 9]

On Page: 303 Line: 9890 Section: exec

In the RATIONALE section

Change From:

"SIG_IGN, and that the process signal mask be unchanged across an exec."

To:

"SIG_IGN, and that the new process image inherits the signal mask of the thread that called exec in the old process image"

Change Number: XSH/TC1/26 [XSH ERN 119]

On Page: 313 Line: 10264 Section: exp

In the EXAMPLES section

Change From:

"None."

To:

"Computing the density of the standard normal distribution

This function shows an implementation for the density of the standard normal distribution using exp(). This example uses the constant M_PI which is an XSI extension.

```
#include <math.h>
```

```
double
```

```
normal_density (double x
```

```
)
```

```
{
```

```
    return exp(-x*x/2) / sqrt (2*M_PI);
```

```
} "
```

Change Number: XSH/TC1/27 [XSH ERN 121]

On Page: 318 Line: 10437 Section: fabs

In the EXAMPLES section

Change From:

"None."

To:

" Computing the 1-norm of a floating point vector

This example shows the use of fabs() to compute the 1-norm of a vector defined as follows:

norm1(v) = |v[0]| + |v[1]| + ... + |v[n-1]|

where |x| denotes the absolute value of x, n denotes the vector's dimension, and v[i] denotes the i-th component of v (0<=i<n)

```
#include <math.h>
```

```
double
```

```
norm1(const double v[], const int n )
```

```

{

    int      i;

    double   n1_v; /* 1-norm of v */

    n1_v = 0;

    for (i=0; i<n; i++) {

        n1_v += fabs (v[i]);

    }

    return n1_v;

}"

```

Change Number: XSH/TC1/28 [XSH ERN 11]

On Page: 327 Line: 10726 Section: fclose

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the process would be delayed in the write operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the thread would be delayed in the write operation. [/CX]"

Rationale: It is the thread that is delayed not the process.

Change Number: XSH/TC1/29 [XSH ERN 134]

On Page: 332 Line: 10935 Section: fcntl

In the EXAMPLES section

Change From:

"None."

To:

"Locking and unlocking a file

The following example demonstrates how to place a lock on bytes 100 to 109 of a file and then later remove it. F_SETLK is used to perform a non-blocking lock request so that the process does not have to wait if an incompatible lock is held by another process; instead the process can take some other action.

```

#include <stdlib.h>

#include <unistd.h>

#include <fcntl.h>

#include <errno.h>

int

main(int argc, char *argv[])

{

    int fd;

    struct flock fl;

    fd = open("testfile", O_RDWR);

```

```

if (fd == -1)

    /* Handle error */;

/* Make a non-blocking request to place a write lock
   on bytes 100-109 of testfile */

fl.l_type = F_WRLCK;

fl.l_whence = SEEK_SET;

fl.l_start = 100;

fl.l_len = 10;

if (fcntl(fd, F_SETLK, &fl) == -1) {

    if (errno == EACCES || errno == EAGAIN) {

        printf("Already locked by another process\n");

        /* We can't get the lock at the moment */

    } else {

        /* Handle unexpected error */;

    }

} else { /* Lock was granted... */

    /* Perform I/O on bytes 100 to 109 of file */

    /* Unlock the locked bytes */

    fl.l_type = F_UNLCK;

    fl.l_whence = SEEK_SET;

    fl.l_start = 100;

    fl.l_len = 10;

    if (fcntl(fd, F_SETLK, &fl) == -1)

        /* Handle error */;

}

exit(EXIT_SUCCESS);

} /* main */

```

Setting the close-on-exec flag

The following example demonstrates how to set the close on exec flag for the file descriptor fd.

```
#include <unistd.h>
```

```
#include <fcntl.h>

...

int flags;

flags = fcntl(fd, F_GETFD);

if (flags == -1)

    /* Handle error */;

flags |= FD_CLOEXEC;

if (fcntl(fd, F_SETFD, flags) == -1)

    /* Handle error */;
```

Change Number: XSH/TC1/30 [XSH ERN 141]

On Page: 343 Line: 11245-11247 Section: fdopen

In the RATIONALE section

Change From:

"The file descriptor may have been obtained from open(), creat(), pipe(), dup(), or fcntl(); inherited through fork() or exec; or perhaps obtained by implementation-defined means, such as the 4.3 BSD socket() call."

To:

The file descriptor may have been obtained from open(), creat(), pipe(), dup(), fcntl() or socket(); inherited through fork(), posix_spawn() or exec; or perhaps obtained by other means."

Change Number: XSH/TC1/31 [XSH ERN 12]

On Page: 360 Line: 11686 Section: fflush

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the process would be delayed in the write operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the thread would be delayed in the write operation. [/CX]"

Change Number: XSH/TC1/32 [XSH ERN 13]

On Page: 364 Line: 11802 Section: fgetc

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the process would be delayed in the fgetc() operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the thread would be delayed in the fgetc() operation. [/CX]"

Change Number: XSH/TC1/33 [XSH ERN 14]

On Page: 370 Line: 11969 Section: fgetwc

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the process would be delayed in the fgetwc() operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the thread would be delayed in the fgetwc() operation. [/CX]"

Change Number: XSH/TC1/34 [XBD ERN 11]

On Page: 399 Line: 12942 Section: fpathconf

In the DESCRIPTION section

Insert after the "{PIPE_BUF}" entry:

"{POSIX2_SYMLINKS} _PC_2_SYMLINKS 4"

Change Number: XSH/TC1/35 [XSH ERN 15]

On Page: 399-402 Line: 12943-12947,12975,13054 Section: fpathconf

In the DESCRIPTION section

Insert the number "10" into the third column for the rows beginning

{POSIX_ALLOC_SIZE_MIN}, {POSIX_REC_INCR_XFER_SIZE},
{POSIX_REC_MAX_XFER_SIZE}, {POSIX_REC_MIN_XFER_SIZE}, and
{POSIX_REC_XFER_ALIGN}.

Add a new bullet item after Item 9 after line 12975

"10. If path or files does not refer to a regular file, it is unspecified whether an implementation supports an association of the variable name with the specified file. If an implementation supports such an association for other than a regular file, the value returned is unspecified."

Add to the RATIONALE section as a new paragraph, after line 13054

"It was the intention of 1003.1d that the variables

{POSIX_ALLOC_SIZE_MIN}, {POSIX_REC_INCR_XFER_SIZE}, {POSIX_REC_MAX_XFER_SIZE},
{POSIX_REC_MIN_XFER_SIZE}, and {POSIX_REC_XFER_ALIGN} only applied to regular files, but note 10
also permits implementation of the advisory semantics on other file types unique to an implementation (e.g. a
character special device.) "

Change Number: XSH/TC1/36 [XBD ERN 27]

On Page: 400 Line: 12991,13020 Section: fpathconf

Add to the end of the RETURN VALUE section as a new paragraph:

"If the variable corresponding to name is dependent on an unsupported option the results are unspecified."

Add to the end of the APPLICATION USAGE section:

"Application writers should check whether an option, such as _POSIX_ADVISORY_INFO is supported, prior to
obtaining and using values for related variables such as POSIX_ALLOC_SIZE_MIN."

Change Number: XSH/TC1/37 [XSH ERN 16]

On Page: 416 Line: 13661 Section: fputc

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the process would
be delayed in the write operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the thread would
be delayed in the write operation. [/CX]"

Change Number: XSH/TC1/38 [XSH ERN 17]

On Page: 420 Line: 13778 Section: fputwc

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the process would
be delayed in the write operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor underlying stream and the thread would
be delayed in the write operation. [/CX]"

Change Number: XSH/TC1/39 [XSH ERN 126]

On Page: 428 Line: 14076 Section: freeaddrinfo

In the RETURN VALUE section

Change From:

"The getaddrinfo() function shall fail and return the corresponding value if:"

To:

"The getaddrinfo() function shall fail and return the corresponding error value if:"

Change Number: XSH/TC1/40 [XSH ERN 142]

On Page: 430 Line: 14141-14144 Section: freopen

In the DESCRIPTION section

Change From:

"If filename is a null pointer, the freopen() function attempt to change the mode of the stream to that by mode, as if the name of the file currently with the stream had been used. It is implementation-defined which changes of mode are permitted (if any) and under what circumstances".

To:

"If filename is a null pointer, the freopen() function shall attempt to change the mode of the stream to that specified by mode, as if the name of the file currently associated with the stream had been used. In this case, the file descriptor associated with the stream need not be closed if the call to freopen() succeeds. It is implementation-defined which changes of mode are permitted (if any) and under what circumstances".

Change Number: XSH/TC1/41 [XSH ERN 142]

On Page: 430,431 Line: 14158,14179 Section: freopen

In the ERRORS section

Add after line 14158

"[CX][EBAF] The file descriptor underlying the stream is not a valid file descriptor when filename is a null pointer. [/CX]"

Add after line 14179

"[CX][EBAF] The mode with which the file descriptor underlying the stream was opened does not support the requested mode when filename is a null pointer. [/CX]"

Change Number: XSH/TC1/42 [XSH ERN 18]

On Page: 443 Line: 14598 Section: fseek

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor and the process would be delayed in the write operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor and the thread would be delayed in the write operation. [/CX]"

Change Number: XSH/TC1/43 [XSH ERN 19]

On Page: 445 Line: 14679 Section: fsetpos

In the ERRORS section

Change From:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor and the process would be delayed in the write operation. [/CX]"

To:

"[CX][EAGAIN] The O_NONBLOCK flag is set for the file descriptor and the thread would be delayed in the write operation. [/CX]"

Change Number: XSH/TC1/44 [XSH ERN 143]

On Page: 452 Line: 14894-14896 Section: fsync

In the DESCRIPTION section

Change From:

"The fsync() function shall request that all data for the open file descriptor named by fildes is to be transferred to the storage device associated with the file described by fildes in an implementation-defined manner."

To:

"The fsync() function shall request that all data for the open file descriptor named by fildes is to be transferred to the storage device associated with the file described by fildes. The nature of the transfer is implementation-defined."

Rationale: This is an editorial wording change with no change intended in meaning.

Change Number: XSH/TC1/45 [XSH ERN 53]

On Page: 491 Line: 16190,16229 Section: getcontext

In the SYNOPSIS section

Change the margin marker from "XSI" to "OB XSI", to denote that the getcontext() and setcontext() functions are obsolescent.

In the APPLICATION USAGE section for getcontext add at the end:

"The obsolescent functions getcontext(), makecontext() and swapcontext() can be replaced using POSIX threads functions."

Change Number: XSH/TC1/46 [XSH ERN 20]

On Page: 554,556 Line: 18219,18266 Section: getrlimit

In the DESCRIPTION section

Change From:

"RLIMIT_STACK This is the maximum size of a process stack, in bytes."

To:

"RLIMIT_STACK This is the maximum size of the initial thread's stack, in bytes."

In the RATIONALE section

Change From:

"None."

To:

"It should be noted that RLIMIT_STACK applies "at least" to the stack of the initial thread in the process, and not to the sum of all the stacks in the process, as that would be very limiting unless the value is so big as to provide no value at all with a single thread."

Change Number: XSH/TC1/47 [XSH ERN 21,22,23]

On Page: 563 Line: 18492-18497 Section: getsockopt

Change From:

"If SO_LINGER is set, the system blocks the process during close() until it can transmit the data or until the end of the interval indicated by the l_linger member, whichever comes first. If SO_LINGER is not specified, and close() is issued, the system handles the call in a way that allows the process to continue as quickly as possible."

To:

"If SO_LINGER is set, the system shall block the calling thread during close() until it can transmit the data or until the end of the interval indicated by the l_linger member, whichever comes first. If SO_LINGER is not specified, and close() is issued, the system handles the call in a way that allows the calling thread to continue as quickly as possible."

Change Number: XSH/TC1/48 [XSH ERN 48]

On Page: 580 Line: 19104-19108 Section: gmtime

In the RETURN VALUE section

Change From:

"Upon successful completion, gmtime_r() shall return the address of the structure pointed to by the argument result. If an error is detected, gmtime_r() shall return a null pointer."

To:

"Upon successful completion, gmtime_r() shall return the address of the structure pointed to by the argument result. If an error is detected, gmtime_r() shall return a null pointer and set errno to indicate the error".

In the ERRORS section:

Change From:

"The gmtime() function shall fail if:

[EOVERFLOW] The result cannot be represented."

To:

"The gmtime() and gmtime_r() function shall fail if:

[EOVERFLOW] The result cannot be represented."

with "and gmtime_r()" shaded and margin marked TSF.

Change Number: XSH/TC1/49 [TC2d5 ERN 3]

On Page: 589 Line: 19374 Section: hypot

In the EXAMPLES section

Change From:

"None."

To:

"See atan2() EXAMPLES."

Change Number: XSH/TC1/50 [XSH ERN 24]

On Page: 648 Line: 21400 Section: isunordered

Change From:

"If x or y is NaN, 0 shall be returned."

To:

"If x or y is NaN, 1 shall be returned."

Rationale:

The previous statement is incorrect. If x is NaN, then isunordered() must return 1. Likewise if y is NaN, then isunordered() must return 1.

Change Number: XSH/TC1/51 [XSH ERN 144]

On Page: 669 Line: 22175-22176 Section: kill

In the RATIONALE section

Change From:

"The implementation-defined processes to which a signal cannot be sent may include the scheduler or init."

To:

"The unspecified processes to which a signal cannot be sent may include the scheduler or init."

Rationale: The RATIONALE section conflicted with the normative text in the DESCRIPTION section.

Change Number: XSH/TC1/52 [XSH ERN 135]

On Page: 671 Line: 22533 Section: killpg

In the EXAMPLES section

Change From:

"None."

To:

"Sending a signal to all other members of a process group

The following example shows how the calling process could send a signal to all other members of its process group. To prevent itself from receiving the signal it first makes itself immune to the signal by ignoring it.

```
#include <signal.h>
```

```
#include <unistd.h>
```

```
...
```

```
if (signal(SIGUSR1, SIG_IGN) == SIG_ERR)
```

```
    /* Handle error */;
```

```
if (killpg(getpgrp(), SIGUSR1) == -1)
```

```
    /* Handle error */;"
```

Change Number: XSH/TC1/53 [XSH ERN 78]

On Page: 686 Line: 22699 Section: lio_listio

In the DESCRIPTION section

Add after line 22699

"If the buffer pointed to by list or the aiocb structures pointed to by the elements of the array list become illegal addresses before all asynchronous I/O completed and, if necessary, the notification is sent, then the behavior is undefined. If the buffers pointed to by the aio_buf member of the aiocb structure pointed to by the elements of the array list become illegal addresses prior to the asynchronous I/O associated with that aiocb structure being completed, the behavior is undefined."

Rationale: This text is added for symmetry with the aio_read() and aio_write() functions.

Change Number: XSH/TC1/54 [XSH ERN 79]

On Page: 686 Line: 22719 Section: lio_listio

In the DESCRIPTION Section (after the last paragraph)

Add after line 22719

"If sig->sigev_notify is SIGEV_THREAD and sig->sigev_notify_attributes is a non-NULL pointer and the block pointed to by this pointer becomes an illegal address prior to all asynchronous I/O being completed, then the behavior is undefined."

Rationale:

This text is added to make it explicit that the user is required to keep the structure pointed to by sig->sigev_notify_attributes valid until the last asynchronous operation finished and the notification has been sent.

Change Number: XSH/TC1/55 [XSH ERN 49]

On Page: 702 Line: 23186-23190 Section: localtime

In the RETURN VALUE

Change From:

"Upon successful completion, localtime_r() shall return a pointer to the structure pointed to by the argument result."

To:

"Upon successful completion, localtime_r() shall return a pointer to the structure pointed to by the argument result. If an error is detected, localtime_r() shall return a null pointer and set errno to indicate the error."

In the ERRORS section

Change From:

"The localtime() function shall fail if:

[EOVERFLOW] The result cannot be represented."

To:

"The localtime() and localtime_r() functions shall fail if:

[EOVERFLOW] The result cannot be represented."

with "and localtime_r()" shaded and margin marked TSF.

Change Number: XSH/TC1/56 [XSH ERN 84,85]

On Page: 702,704 Line: 23181,23245 Section: localtime

In the DESCRIPTION section

Add a new paragraph following the last paragraph with TSF and XSI shading:

"If the reentrant version does not set tzname, it shall not set daylight and shall not set timezone."

In the SEE ALSO section

Add "tzset()," before "utime(),".

Rationale:

On systems supporting XSI, the daylight, timezone, and tzname variables should all be set to provide information for the same timezone.

This updates the description of localtime_r() to mention daylight and timezone as well as tzname.

Change Number: XSH/TC1/57 [XSH ERN 53]

On Page: 732,733 Line: 24081-24083,24143-24146 Section: makecontext

In the SYNOPSIS section

Change the "XSI" margin marker to "OB XSI"

Change the function prototype from:

"void makecontext(ucontext_t *ucp, void (*func)(void), int argc, ...);"

To:

"void makecontext(ucontext_t *ucp, void (*func)(), int argc, ...);"

In the APPLICATION USAGE section

Change From:

"None."

To:

"The obsolescent functions getcontext(), makecontext() and swapcontext() can be replaced using POSIX threads functions."

In the RATIONALE section

Change From:

"None."

To:

"With the incorporation of ISO/IEC 9899:1999 (C99) into this specification it was found that the ISO C Standard (subclause 6.11.6) specifies that the use of function declarators with empty parentheses is an obsolescent feature. Therefore, using the function prototype:

void makecontext(ucontext_t *ucp, void (*func)(), int argc, ...);

is making use of an obsolescent feature of ISO C. Therefore, a strictly conforming POSIX application can't use this form. Therefore, use of getcontext(), makecontext() and swapcontext() is marked obsolescent.

There is no way in ISO C to specify a non-obsolescent function prototype indicating that a function will be called with an arbitrary number (including zero) of arguments of arbitrary types (including integers, pointers to data, pointers to functions, and composite types). Replacing makecontext() with a number of ISO C compatible functions handling various numbers and types of arguments would have forced all existing uses of makecontext() to be rewritten for little or no gain.

There are very few applications today that use the *context() routines. Those that do use them are almost always using them to implement co-routines. By maintaining the XSH5 specification for makecontext(), existing applications will continue to work, although they won't be able to be classified as strictly conforming applications.

There is no way in ISO C (without using obsolescent behavior) to specify functionality that was standard, strictly conforming behavior in XSH5 using the 1990 C Standard. Threads can be used to implement the functionality provided by makecontext(), getcontext() and swapcontext() but they are more complex to use. It was felt inventing new ISO C compatible interfaces that describe what can be done with the XSH5 functions and then converting applications to use them would cause more difficulty than just converting applications that use them to use threads instead."

Change Number: XSH/TC1/58 [XSH ERN 50]

On Page: 768 Line: 25150-25152,25154 Section: mktime

In the RETURN VALUE section

Change From:

"The mktime() function shall return the specified time since the Epoch encoded as a value of type time_t. If the time since the Epoch cannot be represented, the function shall return the value (time_t)-1."

To:

"The mktime() function shall return the specified time since the Epoch encoded as a value of type time_t. If the time since the Epoch cannot be represented, the function shall return the value (time_t)-1 [CX] and may set errno to indicate the error.[CX]."

In the ERRORS section

Change From:

"None."

To:

(Shaded and marked with the CX margin code):

"The mktime() function may fail if:

[EOVERFLOW] The result cannot be represented."

Change Number: XSH/TC1/59 [XSH ERN 86]

On Page: 769 Line: 25185 Section: mktime

In the SEE ALSO section

Add "tzset()," before "utime(),"

Rationale: This is editorial.

Change Number: XSH/TC1/60 [XSH ERN 72]

On Page: 774,777 Line: 25357,25469 Section: mmap

In the DESCRIPTION section

Add a new paragraph after line 25357

"If len is zero mmap() shall fail and no mapping shall be established."

In the ERRORS section (shall fail)

Add after line 25469

"[EINVAL] The value of len is zero."

Change Number: XSH/TC1/61 [XSH ERN 47]

On Page: 786 Line: 25817 Section: mq_getattr

In the ERRORS section

Change From:

"The mq_getattr() function shall fail if:"

To:

"The mq_getattr() function may fail if:"

Change Number: XSH/TC1/62 [XSH ERN 87]

On Page: 791 Line: 25933 Section: mq_open

In the DESCRIPTION section (O_CREAT mode)

Change From:

"The file permission bits shall be set to the value of mode. When bits in mode other than file permission bits are set, the effect is implementation-defined."

To:

"The permission bits of the message queue shall be set to the value of the mode argument except those set in the file mode creation mask of the process. When bits in mode other than the file permission bits are specified, the effect is unspecified."

Rationale: Consistency with shm_open(), sem_open().

Change Number: XSH/TC1/63 [XSH ERN 25]

On Page: 823 Line: 26879-26880 Section: nanosleep

In the RATIONALE section

Change From:

"It is common to suspend execution of a process for an interval in order to poll the status of a non-interrupting function."

To:

"It is common to suspend execution of a thread for an interval in order to poll the status of a non-interrupting function."

Change Number: XSH/TC1/64 [XSH ERN 73]

On Page: 829-831 Line: 27026,27069,27102,27114 Section: nftw

In the SYNOPSIS section, line 27026

Change From:

"int nftw(const char *path, int (*fn)(const char *, const struct stat *,
int, struct FTW *), int depth, int flags);"

To:

"int nftw(const char *path, int (*fn)(const char *, const struct
stat *, int, struct FTW *), int fd_limit, int flags);"

In the DESCRIPTION section, line 27069

Change From:

"The argument depth sets the maximum number of file descriptors that shall be used by nftw() while traversing the file tree."

To:

"The argument `fd_limit` sets the maximum number of file descriptors that shall be used by `nftw()` while traversing the file tree."

In the EXAMPLES section, line 27102

Change From:

"The following example walks the `/tmp` directory and its subdirectories, calling the `nftw()` function for every directory entry, to a maximum of 5 levels deep."

To:

"The following example walks the `/tmp` directory and its subdirectories, calling the `nftw()` function for every directory entry, using a maximum of 5 file descriptors."

On line 27114

Change From:

```
"int depth = 5;
```

```
int flags = FTW_CHDIR | FTW_DEPTH | FTW_MOUNT;
```

```
int ret;
```

```
ret = nftw(startpath, nftwfunc, depth, flags);"
```

To:

```
"int fd_limit = 5;
```

```
int flags = FTW_CHDIR | FTW_DEPTH | FTW_MOUNT;
```

```
int ret;
```

```
ret = nftw(startpath, nftwfunc, fd_limit, flags);"
```

Change Number: XSH/TC1/65 [XSH ERN 127]

On Page: 856 Line: 27866 Section: pipe

In the EXAMPLES section

Change From:

"None."

To:

"Using a pipe to pass data between a parent process and a child process

The following example demonstrates the use of a pipe to transfer data between a parent process and a child process. Error handling is excluded, but otherwise this code demonstrates good practice when using pipes: after the `fork()` the two processes close the unused ends of the pipe before they commence transferring data.

```
#include <stdlib.h>
```

```
#include <unistd.h>
```

```
...
```

```
int fildes[2];
```

```
const int BSIZE = 100;
```

```
char buf[BSIZE];
```

```
ssize_t nbytes;
```

```
int status;
```

```
status = pipe(fildes);
```

```

if (status == -1 ) {

    /* an error occurred */

    ....

}

switch (fork()) {

case -1: /* Handle error */

    break;

case 0: /* Child - reads from pipe */

    close(fildes[1]); /* Write end is unused */

    nbytes = read(fildes[0], buf, BSIZE); /* Get data from pipe */

    /* At this point, a further read would see end of file... */

    close(fildes[0]); /* Finished with pipe */

    exit(EXIT_SUCCESS);

default: /* Parent - writes to pipe */

    close(fildes[0]); /* Read end is unused */

    write(fildes[1], "Hello world\n", 12); /* Write data on pipe */

    close(fildes[1]); /* Child will see EOF */

    exit(EXIT_SUCCESS);

}"

```

Change Number: XSH/TC1/66 [XSH ERN 145]

On Page: 860 Line: 27995-27999 Section: poll

In the EXAMPLES section

Change From:

```

"/* Open STREAMS device. */

fds[0].fd = open("/dev/dev0", ...);

fds[1].fd = open("/dev/dev1", ...);

fds[0].events = POLLOUT | POLLWRBAND;

fds[1].events = POLLOUT | POLLWRBAND;"

```

To:

```

/* Open STREAMS device. */

fds[0].fd = open("/dev/dev0", ...);

fds[1].fd = open("/dev/dev1", ...);

fds[0].events = POLLOUT | POLLWRBAND;

fds[1].events = POLLOUT | POLLWRBAND;"

```

Change Number: XSH/TC1/67 [XSH ERN 128]

On Page: 863 Line: 28077 Section: popen

In the EXAMPLES section

Change From:

"None."

To:

"Using popen() to obtain a list of files from the ls utility

The following example demonstrates the use of popen() and pclose() to execute the command "ls *" in order to obtain a list of files in the current directory:

```

#include <stdio.h>

...

FILE *fp;

int status;

char path[PATH_MAX];

fp = popen("ls *", "r");

if (fp == NULL)

    /* Handle error */;

while (fgets(path, PATH_MAX, fp) != NULL)

    printf("%s", path);

status = pclose(fp);

if (status == -1) {

    /* Error reported by pclose() */

    ...

} else {

    /* Use macros described under wait() to inspect 'status' in order
       to determine success/failure of command executed by popen() */

    ...

}"

```

Change Number: XSH/TC1/68 [XSH ERN 146]

On Page: 864 Line: 28122 Section: posix_fadvise

In the SYNOPSIS section

Change From:

"int posix_fadvise(int fd, off_t offset, size_t len, int advice);"

To:

"int posix_fadvise(int fd, off_t offset, off_t len, int advice);"

Rationale: The previous prototype was not large-file aware, and the standard developers felt it acceptable to make this change before implementations of the functions become widespread."

Change Number: XSH/TC1/69 [XSH ERN 146]

On Page: 866 Line: 28177 Section: posix_fallocate

In the SYNOPSIS section

Change From:

"int posix_fallocate(int fd, off_t offset, size_t len);"

To:

"int posix_fallocate(int fd, off_t offset, off_t len);"

Rationale: The previous prototype was not large-file aware, and the standard developers felt it acceptable to make this change before implementations of the functions become widespread."

Change Number: XSH/TC1/70 [XSH ERN 110]

On Page: 975 Line: 31069-31071 Section: pselect

In the DESCRIPTION section

Change From:

"If sigmask is not a null pointer, then the pselect() function shall replace the signal mask of the process by the set of signals pointed to by sigmask before examining the descriptors, and shall restore the signal mask of the process before returning."

To:

"If sigmask is not a null pointer, then the pselect() function shall replace the signal mask of the caller by the set of signals pointed to by sigmask before examining the descriptors, and shall restore the signal mask of the calling thread before returning."

Change Number: XSH/TC1/71 [XSH ERN 147,148]

On Page: 981 Line: 31293 Section: pthread_attr_destroy

In the ERRORS section

Add

"The pthread_attr_destroy() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object.

The pthread_attr_init() may fail if:

[EBUSY] The implementation has detected an attempt to reinitialize the thread attribute referenced by attr, a previously initialized, but not yet destroyed, thread attribute."

Change Number: XSH/TC1/72 [XSH ERN 125]

On Page: 984 Line: 31407 Section: pthread_attr_getdetachstate

In the EXAMPLES section

Change From:

"None."

To:

"Retrieving the detachstate attribute.

This example shows how to obtain the detachstate attribute of a thread attribute object.

```
#include <pthread.h>
```

```
pthread_attr_t thread_attr;
```

```
int detachstate;
```

```
int rc;
```



```

/* code initializing thread_attr */

...

rc = pthread_attr_getdetachstate (&thread_attr, &detachstate);

if (rc!=0) {

    /* handle error */

    ...

}

else {

    /* legal values for detachstate are:

    * PTHREAD_CREATE_DETACHED or PTHREAD_CREATE_JOINABLE

    */

    ...

} "

```

Change Number: XSH/TC1/73 [XSH ERN 149]

On Page: 984 Line: 31404 Section: pthread_attr_getdetachstate

In the ERRORS section

Add

"The pthread_attr_getdetachstate() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object.

The pthread_attr_setdetachstate() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

Change Number: XSH/TC1/74 [XSH ERN 150]

On Page: 986 Line: 31461,31462 Section: pthread_attr_getdetachstate

In the ERRORS section

On line 31461 delete:

"[EINVAL] The attribute attr is invalid."

In the ERRORS section

Add:

"The pthread_attr_getguardsize() and pthread_attr_setguardsize may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

Change Number: XSH/TC1/75 [XSH ERN 151,152]

On Page: 986 Line: 31500,31516,31517 Section: pthread_attr_getinheritsched

In the DESCRIPTION section

After line 31500 insert a new paragraph:

"The supported values of inheritsched shall be:"

In the ERRORS section

Insert before line 31516

"The pthread_attr_getinheritsched() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

In the ERRORS section

Add after line 31517 :

"[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

Change Number: XSH/TC1/76 [XSH ERN 124]

On Page: 987 Line: 31465 Section: pthread_attr_getguardsize

In the EXAMPLES section

Change From:

"None."

To:

"Retrieving the guardsize attribute."

This example shows how to obtain the guardsize attribute of a thread attribute object.

```
#include <pthread.h>

pthread_attr_t thread_attr;

size_t  guardsize;

int      rc;

/* code initializing thread_attr */

...

rc = pthread_attr_getguardsize (&thread_attr, &guardsize);

if (rc != 0) {

    /* handle error */

    ...

}

else {

    if (guardsize > 0) {

        /* a guard area of at least guardsize bytes is provided */

        ...

    }

    else {

        /* no guard area provided */

        ...

    }

}"
```

Change Number: XSH/TC1/77 [XSH ERN 55]

On Page: 988 Line: 31524 Section: pthread_attr_getinheritsched

In the APPLICATION USAGE section

Add:

"See Section 2.9.4 for further details on thread scheduling attributes and their default settings."

Change Number: XSH/TC1/78 [XSH ERN 154]

On Page: 990 Line: 31567,31569 Section: pthread_attr_getschedparam

In the ERRORS section

Add after line 31567 :

"The pthread_attr_getschedparam() function may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

On line 31569

Change From:

"[EINVAL] The value of param is not valid."

To:

"[EINVAL] The value of param is not valid, or the value specified by attr does not refer to an initialized thread attribute object."

Change Number: XSH/TC1/79 [XSH ERN 56]

On Page: 992 Line: 31621 Section: pthread_attr_getschedpolicy

In the APPLICATION USAGE section

Add:

"See Section 2.9.4 for further details on thread scheduling attributes and their default settings."

Change Number: XSH/TC1/80 [XSH ERN 155]

On Page: 992 Line: 31612,31614 Section: pthread_attr_getschedpolicy

In the ERRORS section

Add after line 31612 :

"The pthread_attr_getschedpolicy() function may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

On line 31614

Change From:

"[EINVAL] The value of policy is not valid."

To:

"[EINVAL] The value of policy is not valid, or the value specified by attr does not refer to an initialized thread attribute object."

Change Number: XSH/TC1/81 [XSH ERN 57]

On Page: 994 Line: 31669 Section: pthread_attr_getscope

In the APPLICATION USAGE section

Add:

"See Section 2.9.4 for further details on thread scheduling attributes and their default settings."

Change Number: XSH/TC1/82 [XSH ERN 157]

On Page: 994 Line: 31660,31662 Section: pthread_attr_getscope

In the ERRORS section

Add after line 31660 :

"The pthread_attr_getscope() function may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

On line 31662

Change From:

"[EINVAL] The value of contentionscope is not valid."

To:

"[EINVAL] The value of contentionscope is not valid, or the value specified by attr does not refer to an initialized thread attribute object."

Change Number: XSH/TC1/83 [XSH ERN 89]

On Page: 996 Line: 31731 Section: pthread_attr_getstack

In the APPLICATION USAGE

Add at the end:

"After a successful call to pthread_attr_setstack(), the storage area specified by the stackaddr parameter is under the control of the implementation as described in Section 2.9.8."

Change Number: XSH/TC1/84 [XSH ERN 158]

On Page: 996 Line: 31714,31716 Section: pthread_attr_getstack

In the ERRORS section

Add after line 31774 :

"The pthread_attr_getstack() function may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

On line 31716

Change From:

"[EINVAL] The value of stackaddr does not have proper alignment to be used as a stack, or if (stackaddr + stacksize) lacks proper alignment."

To:

"[EINVAL] The value of stackaddr does not have proper alignment to be used as a stack, or (stackaddr + stacksize) lacks proper alignment, or the value specified by attr does not refer to an initialized thread attribute object."

Change Number: XSH/TC1/85 [XSH ERN 89]

On Page: 998 Line: 31777 Section: pthread_attr_getstackaddr

In the APPLICATION USAGE

Add at the end:

"After a successful call to pthread_attr_setstackaddr(), the storage area specified by the stackaddr parameter is under the control of the implementation as described in Section 2.9.8."

Change Number: XSH/TC1/86 [XSH ERN 159]

On Page: 998 Line: 31761 Section: pthread_attr_getstackaddr

In the ERRORS section

Change From:

"No errors are defined."

To:

"The pthread_attr_getstackaddr() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object.

The pthread_attr_setstackaddr() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

Change Number: XSH/TC1/87 [XSH ERN 160]

On Page: 1000 Line: 31815 Section: pthread_attr_getstacksize

In the ERRORS section

Add:

"The pthread_attr_getstacksize() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object.

The pthread_attr_setstacksize() may fail if:

[EINVAL] The value specified by attr does not refer to an initialised thread attribute object."

Change Number: XSH/TC1/88 [XSH ERN 64]

On Page: 1024 Line: 32233 Section: pthread_cleanup_pop

In the DESCRIPTION section

Add a new final paragraph:

"The effect of the use of 'return', 'break', 'continue', and 'goto' to prematurely leave a code block described by a pair of pthread_cleanup_push() and pthread_cleanup_pop() functions calls is undefined."

Change Number: XSH/TC1/89 [XSH ERN 90]

On Page: 1036 Line: 32652-32655 Section: pthread_cond_timedwait

In the DESCRIPTION section

Change From:

"The effect of using more than one mutex for concurrent pthread_cond_timedwait() or pthread_cond_wait() operations on the same condition variable is undefined; that is, a condition variable becomes bound to a unique mutex when a thread waits on the condition variable, and this (dynamic) binding shall end when the wait returns."

To:

"When a thread waits on a condition variable, having specified a particular mutex to either the `pthread_cond_timedwait()` or the `pthread_cond_wait()` operation, a dynamic binding is formed between that mutex and condition variable that remains in effect as long as at least one thread is blocked on the condition variable. During this time, the effect of an attempt by any thread to wait on that condition variable using a different mutex is undefined. Once all waiting threads have been unblocked (as by the `pthread_cond_broadcast()` operation), the next wait operation on that condition variable shall form a new dynamic binding with the mutex specified by that wait operation. Even though the dynamic binding between condition variable and mutex may be removed or replaced between the time a thread is unblocked from a wait on the condition variable and the time that it returns to the caller or begins cancellation cleanup, the unblocked thread shall always re-acquire the mutex specified in the condition wait operation call from which it is returning."

Rationale:

This change is for consistency with `pthread_cond_destroy()` which says it is safe to destroy an initialized condition variable upon which no threads are currently blocked.

Change Number: XSH/TC1/90 [XSH ERN 58]

On Page: 1036 Line: 32656-32659 Section: `pthread_cond_timedwait`

In the DESCRIPTION section

Change From:

"When the cancelability enable state of a thread is set to `PTHREAD_CANCEL_DEFERRED`, a side effect of acting upon a cancellation request while in a condition wait is that the mutex is (in effect) re-acquired before calling the first cancellation cleanup handler."

To:

"When the cancelability type of a thread is set to `PTHREAD_CANCEL_DEFERRED`, a side effect of acting upon a cancellation request while in a condition wait is that the mutex is (in effect) re-acquired before calling the first cancellation cleanup handler."

Change Number: XSH/TC1/91 [XSH ERN 26]

On Page: 1037 Line: 32690-32693 Section: `pthread_cond_timedwait`

In the ERRORS section

Change From:

"The `pthread_cond_timedwait()` and `pthread_cond_wait()` functions may fail if:

[EINVAL] The value specified by `cond`, `mutex`, or `abstime` is invalid.

[EINVAL] Different mutexes were supplied for concurrent `pthread_cond_timedwait()` or `pthread_cond_wait()` operations on the same condition variable."

To:

"The `pthread_cond_timedwait()` function shall fail if:

[ETIMEDOUT] The time specified by `abstime` to `pthread_cond_timedwait()` has passed.

[EINVAL] The value specified by `abstime` is invalid.

The `pthread_cond_timedwait()` and `pthread_cond_wait()` functions may fail if:

[EINVAL] The value specified by `cond` or `mutex` is invalid."

Change Number: XSH/TC1/92 [XSH ERN 118]

On Page: 1037 Line: 32705 Section: `pthread_cond_timedwait`

In the RATIONALE section

Add a new second paragraph:

"The application needs to recheck the predicate on any return because it cannot be sure there is another thread waiting on the thread to handle the signal, and if there is not then the signal is lost. The burden is on the application to check the predicate."

Change Number: XSH/TC1/93 [XSH ERN 136]

On Page: 1050 Line: 33036,33038 Section: `pthread_create`

In the ERROR RETURN section

Delete the following from the "shall fail" section:

"[EINVAL] The value specified by `attr` is invalid."

Add after line 33038 :

"The `pthread_create()` function may fail if:

[EINVAL] The attributes specified by `attr` are invalid."

Change Number: XSH/TC1/94 [XSH ERN 91]

On Page: 1051 Line: 33043 Section: pthread_create

In the APPLICATION USAGE section

Change From:

"None."

To:
"There is no requirement on the implementation that the ID of the created thread be available before the newly created thread starts executing. The calling thread can obtain the ID of the created thread through the return value of the pthread_create() function, and the newly created thread can obtain its ID by a call to pthread_self()."

Change Number: XSH/TC1/95 [XSH ERN 27]

On Page: 1053 Line: 33123-33127 Section: pthread_detach

In the ERRORS section

Change From:

"The pthread_detach() function shall fail if:

[EINVAL] The implementation has detected that the value specified by thread does not refer to a joinable thread.

[ESRCH] No thread could be found corresponding to that specified by the given thread ID."

To:
"The pthread_detach() function may fail if:
[EINVAL] The implementation has detected that the value specified by thread does not refer to a joinable thread.
[ESRCH] No thread could be found corresponding to that specified by the given thread ID."

Change Number: XSH/TC1/96 [XSH ERN 100]

On Page: 1064 Line: 33451 Section: pthread_getspecific

In the ERRORS Section

Change From:

"[ENOMEM] Insufficient memory exists to associate the value with the key."

To:
"[ENOMEM] Insufficient memory exists to associate the non-NULL value with the key."

Change Number: XSH/TC1/97 [XSH ERN 28]

On Page: 1066 Line: 33494-33500 Section: pthread_join

Change From:

"The pthread_join() function shall fail if:

[EINVAL] The implementation has detected that the value specified by thread does not refer to a joinable thread.

[ESRCH] No thread could be found corresponding to that specified by the given thread ID.

The pthread_join() function may fail if:

[EDEADLK] A deadlock was detected or the value of thread specifies the calling thread."

To:
"The pthread_join() function shall fail if:
[ESRCH] No thread could be found corresponding to that specified by the given thread ID.
The pthread_join() function may fail if:
[EDEADLK] A deadlock was detected or the value of thread specifies the calling thread.
[EINVAL] The value specified by thread does not refer to a joinable thread."

Change Number: XSH/TC1/98 [XSH ERN 59,65]

On Page: 1085 Line: 34135,34150 Section: pthread_mutex_lock

In the ERRORS section

Change From:

"[EDEADLK] The current thread already owns the mutex."

To:
"[EDEADLK] A deadlock condition was detected or the current thread already owns the mutex."

In the RATIONALE section

Change From:

"For example, deadlocking on a double-lock is explicitly allowed behavior in order to avoid requiring more overhead in the basic mechanism than is absolutely necessary."

To:

"For example, on systems not supporting the XSI extended mutex types, deadlocking on a double-lock is explicitly allowed behavior in order to avoid requiring more overhead in the basic mechanism than is absolutely necessary."

Add to the end of the RATIONALE section:

"For further rationale on the XSI extended mutex types please see the Rationale Volume."

Change Number: XSH/TC1/99 [XSH ERN 60]

On Page: 1088 Line: 34211-34214 Section: pthread_mutex_timedlock

In the DESCRIPTION section (last paragraph)

Margin mark the following text TPI and shade:

"As a consequence of the priority inheritance rules (for mutexes initialized with the PRIO_INHERIT protocol), if a timed mutex wait is terminated because its timeout expires, the priority of the owner of the mutex shall be adjusted as necessary to reflect the fact that this thread is no longer among the threads waiting for the mutex."

Change Number: XSH/TC1/100 [XSH ERN 65]

On Page: 1089 Line: 34231 Section: pthread_mutex_timedlock

In the ERRORS section

Change From:

"[EDEADLK] The current thread already owns the mutex."

To:

"[EDEADLK] A deadlock condition was detected or the current thread already owns the mutex."

Change Number: XSH/TC1/101 [XSH ERN 65]

On Page: 1114 Line: 34988 Section: pthread_rwlock_rdlock

In the ERRORS section

Change From:

"[EDEADLK] The current thread already owns the read-write lock for writing."

To:

"[EDEADLK] A deadlock condition was detected or the current thread already owns the read-write lock for writing."

Change Number: XSH/TC1/102 [XSH ERN 65]

On Page: 1115 Line: 35054 Section: pthread_rwlock_timedrdlock

In the ERRORS section

Change From:

"[EDEADLK] The calling thread already holds a write lock on rwlock."

To:

"[EDEADLK] A deadlock condition was detected or the calling thread already holds a write lock on rwlock."

Change Number: XSH/TC1/103 [XSH ERN 65]

On Page: 1117 Line: 35112 Section: pthread_rwlock_timedwrlock

In the ERRORS section

Change From:

"[EDEADLK] The calling thread already holds the rwlock."

To:

"[EDEADLK] A deadlock condition was detected or the calling thread already holds the rwlock."

Change Number: XSH/TC1/104 [XSH ERN 65]

On Page: 1120 Line: 35178 Section: pthread_rwlock_trywrlock

In the ERRORS section

Change From:

"[EDEADLK] The current thread already owns the read-write lock for writing or reading."

To:

"[EDEADLK] A deadlock condition was detected or the current thread already owns the read-write lock for writing or reading."

Change Number: XSH/TC1/105 [XSH ERN 30]

On Page: 1139-1140 Line: 35589,35615 Section: pthread_sigmask

In the DESCRIPTION section, line 35589

Change From:

"If set is a null pointer, the value of the argument how is not significant and the process signal mask shall be unchanged; thus the call can be used to enquire about currently blocked signals."

To:

"If set is a null pointer, the value of the argument how is not significant and the thread signal mask shall be unchanged; thus the call can be used to enquire about currently blocked signals."

In the RATIONALE section, line 35615

Change From:

"When a process signal mask is changed in a signal-catching function that is installed by sigaction(), the restoration of the signal mask on return from the signal-catching function overrides that change (see sigaction())."

To:

"When a thread's signal mask is changed in a signal-catching function that is installed by sigaction(), the restoration of the signal mask on return from the signal-catching function overrides that change (see sigaction())."

Change Number: XSH/TC1/106 [XSH ERN 114]

On Page: 1140 Line: 35611 Section: pthread_sigmask

In the EXAMPLES section

Change From:

"None."

To:

"Signalling in multi-threaded process."

This example shows the use of pthread_sigmask() in order to deal with signals in a multi-threaded process. It provides a fairly general framework that could be easily adapted/extended by the reader.

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
#include <pthread.h>
```

```
#include <signal.h>
```

```
#include <string.h>
```

```
#include <errno.h>
```

```
...
```

```
static sigset_t signal_mask; /* signals to block */
```

```
int main (int argc, char *argv[])
```

```
{
```

```
pthread_t sig_thr_id; /* signal handler thread ID */
```

```
int rc; /* return code */
```

```
sigemptyset (&signal_mask);
```

```
sigaddset (&signal_mask, SIGINT);
```

```
sigaddset (&signal_mask, SIGTERM);
```

```
rc = pthread_sigmask (SIG_BLOCK, &signal_mask, NULL);
```



```

if (rc != 0) {
    /* handle error */
    ...
}

/* any newly created threads inherit the signal mask */
rc = pthread_create (&sig_thr_id, NULL, signal_thread, NULL);
if (rc != 0) {
    /* handle error */
    ...
}

/* APPLICATION CODE */
...
}

void *signal_thread (void *arg)
{
    int      sig_caught;    /* signal caught      */
    int      rc;            /* returned code   */

    rc = sigwait (&signal_mask, &sig_caught);
    if (rc != 0) {
        /* handle error */
    }

    switch (sig_caught)
    {
    case SIGINT:    /* process SIGINT */
        ...
        break;

    case SIGTERM:   /* process SIGTERM */
        ...
        break;

    default:        /* should normally not happen */

```

```
        fprintf (stderr, "\nUnexpected signal %d\n", sig_caught);

        break;

    }

}
```

Change Number: XSH/TC1/107 [XSH ERN 65]

On Page: 1143 Line: 35723 Section: pthread_spin_lock

In the ERRORS section

Change From:

"[EDEADLK] The calling thread already holds the lock."

To:

"[EDEADLK] A deadlock condition was detected or the calling thread already holds the lock."

Change Number: XSH/TC1/108 [XSH ERN 43]

On Page: 1174 Line: 36533 Section: read

In the ERRORS section

Change From:

"[EAGAIN] The O_NONBLOCK flag is set for the file descriptor and the process would be delayed."

To:

"[EAGAIN] The O_NONBLOCK flag is set for the file descriptor and the thread would be delayed."

Change Number: XSH/TC1/109 [XSH ERN 165]

On Page: 1177 Line: 36652-36654 Section: read

In the RATIONALE section

Change From:

"IEEE Std 1003.1-2001 does not specify when an implementation that buffers read()ss actually moves the data into the user-supplied buffer, so an implementation may chose to do this at the latest possible moment."

To:

"IEEE Std 1003.1-2001 does not specify when an implementation that buffers read()s actually moves the data into the user-supplied buffer, so an implementation may chose to do this at the latest possible moment."

Change Number: XSH/TC1/110 [XSH ERN 111]

On Page: 1190 Line: 37082,37091 Section: realpath()

Change From:

"[ELOOP] A loop exists in symbolic links encountered during resolution of the path argument."

To:

"[ELOOP] A loop exists in symbolic links encountered during resolution of the file_name argument."

Change From:

"[ELOOP] More than {SYMLOOP_MAX} symbolic links were encountered during resolution of the path argument."

To:

"[ELOOP] More than {SYMLOOP_MAX} symbolic links were encountered during resolution of the file_name argument."

Rationale:

These ERROR descriptions erroneously referred to a non-existent path argument.

Change Number: XSH/TC1/111 [XSH Errata 2003-1]

On Page: 1247 Line: 38887 Section: select

In the SYNOPSIS section

Change from:

"#include <sys/time.h>"

To:

"#include <sys/select.h>"

Rationale: This is an editorial correction to the pointer page to match the actual manual page.

Change Number: XSH/TC1/112 [XSH ERN 66]

On Page: 1248 Line: 38915 Section: sem_close

In the ERRORS section

Change From:

"The sem_close() function shall fail if:"

To:

"The sem_close() function may fail if:"

Change Number: XSH/TC1/113 [XSH ERN 67]

On Page: 1249 Line: 38952-38955 Section: sem_destroy

In the ERRORS section

Change From:

"The sem_destroy() function shall fail if:

[EINVAL] The sem argument is not a valid semaphore.

The sem_destroy() function may fail if:

[EBUSY] There are currently processes blocked on the semaphore."

To:

"The sem_destroy() function may fail if:

[EINVAL] The sem argument is not a valid semaphore.

[EBUSY] There are currently processes blocked on the semaphore."

Change Number: XSH/TC1/114 [XSH ERN 68]

On Page: 1250 Line: 38993 Section: sem_getvalue

In the ERRORS section

Change From:

"The sem_getvalue() function shall fail if:"

To:

"The sem_getvalue() function may fail if:"

Change Number: XSH/TC1/115 [XSH ERN 102]

On Page: 1252 Line: 39025-39036 Section: sem_init

In the DESCRIPTION section

Change From:

"The sem_init() function shall initialize the unnamed semaphore referred to by sem. The value of the initialized semaphore shall be value. Following a successful call to sem_init(), the semaphore may be used in subsequent calls to sem_wait(), sem_trywait(), sem_post(), and sem_destroy(). This semaphore shall remain usable until the semaphore is destroyed.

If the pshared argument has a non-zero value, then the semaphore is shared between processes; in this case, any process that can access the semaphore sem can use sem for performing sem_wait(), sem_trywait(), sem_post(), and sem_destroy() operations.

Only sem itself may be used for performing synchronization. The result of referring to copies of sem in calls to sem_wait(), sem_trywait(), sem_post(), and sem_destroy() is undefined.

If the pshared argument is zero, then the semaphore is shared between threads of the process; any thread in this process can use sem for performing sem_wait(), sem_trywait(), sem_post(), and sem_destroy() operations."

To:

"The sem_init() function shall initialize the unnamed semaphore referred to by sem. The value of the initialized semaphore shall be value. Following a successful call to sem_init(), the semaphore may be used in subsequent calls to sem_wait(), [TMO]sem_timedwait(), [TMO] sem_trywait(), sem_post(), and sem_destroy(). This semaphore shall remain usable until the semaphore is destroyed.

If the pshared argument has a non-zero value, then the semaphore is shared between processes; in this case, any process that can access the semaphore sem can use sem for performing sem_wait(), [TMO]sem_timedwait(), [TMO]sem_trywait(), sem_post(), and sem_destroy() operations.

Only sem itself may be used for performing synchronization. The result of referring to copies of sem in calls to sem_wait(), [TMO]sem_timedwait(), [TMO]sem_trywait(), sem_post(), and sem_destroy() is undefined.

If the pshared argument is zero, then the semaphore is shared between threads of the process; any thread in this process can use sem for performing sem_wait(), [TMO]sem_timedwait(), [TMO]sem_trywait(), sem_post(), and sem_destroy() operations."

Rationale: The sem_timedwait() function is added to the DESCRIPTION for alignment with 1003.1d-1999.

Change Number: XSH/TC1/116 [XSH ERN 103]

On Page: 1254 Line: 39079-39080 Section: sem_open

In the DESCRIPTION section

Change From:

"This semaphore may be used in subsequent calls to sem_wait(), sem_trywait(), sem_post(), and sem_close()."

To:

"This semaphore may be used in subsequent calls to sem_wait(), [TMO]sem_timedwait(),[TMO] sem_trywait(), sem_post(), and sem_close()."

Rationale: The sem_timedwait() function is added to the DESCRIPTION for alignment with 1003.1d-1999.

Change Number: XSH/TC1/117 [XSH ERN 97]

On Page: 1254-1255 Line: 39115-39117 Section: sem_open

In the DESCRIPTION section

Change From:

"If a process makes multiple successful calls to sem_open() with the same value for name, the same semaphore address shall be returned for each such successful call, provided that there have been no calls to sem_unlink() for this semaphore."

To:

"If a process makes multiple successful calls to sem_open() with the same value for name, the same semaphore address shall be returned for each such successful call, provided that there have been no calls to sem_unlink() for this semaphore and at least one previous successful sem_open() call for this semaphore has not been matched with sem_close() call."

Change Number: XSH/TC1/118 [XSH ERN 69]

On Page: 1257 Line: 39192 Section: sem_post

In the ERRORS section

Change From:

"The sem_post() function shall fail if:"

To:

"The sem_post() function may fail if:"

Change Number: XSH/TC1/119 [XSH ERN 70]

On Page: 1259 Line: 39247 Section: sem_timedwait

In the ERRORS section

Move the following text on line 39247 from the "shall fail" to the "may fail" section:

"[EINVAL] The sem argument does not refer to a valid semaphore."

Change Number: XSH/TC1/120 [XSH ERN 71]

On Page: 1261 Line: 39298 Section: sem_trywait

In the ERRORS section

Move the following text on line 39298 from the "shall fail" to the "may fail" section:

"[EINVAL] The sem argument does not refer to a valid semaphore."

Change Number: XSH/TC1/121 [XSH ERN 133]

On Page: 1269 Line: 39510 Section: semget

In the DESCRIPTION section

Change From:

"The data structure associated with each semaphore in the set shall not be initialized."

To:

"The data structure associated with each semaphore in the set need not be initialized."

Change Number: XSH/TC1/122 [XSH ERN 44]

On Page: 1290 Line: 40250-40251 Section: seteuid

In the RATIONALE section

Change From:

"[EPERM] The process does not have appropriate privileges and uid does not match the real group ID or the saved set-group-ID."

To:

"[EPERM] The process does not have appropriate privileges and uid does not match the real user ID or the saved set-user-ID."

Rationale. This is a correction of an editorial error.

Change Number: XSH/TC1/123 [XSH ERN 92]

On Page: 1299 Line: 40449,40456 Section: setlocale

In the DESCRIPTION section

Change From:

" "" Specifies an implementation-defined native environment. [CX]

This corresponds to the value of the associated environment variables, LC_* and LANG ; see the Base Definitions volume of IEEE Std 1003.1-2001, Chapter 7, Locale and the Base Definitions volume of IEEE Std 1003.1-2001, Chapter 8, Environment Variables. [/CX] "

To:

" "" Specifies an implementation-defined native environment. [CX]

The determination of the name of the new locale for the specified category depends on the value of the associated environment variables, LC_* and LANG ; see the Base Definitions volume of IEEE Std 1003.1-2001, Chapter 7, Locale and the Base Definitions volume of IEEE Std 1003.1-2001, Chapter 8, Environment Variables. [/CX] "

In the DESCRIPTION section

Insert before the last paragraph:

"[CX] Setting all of the categories of the locale of the process is similar to successively setting each individual category of the locale of the process, except that all error checking is done before any actions are performed.

To set all the categories of the locale of the process, setlocale() is invoked as:

setlocale(LC_ALL, "");

In this case, setlocale() shall first verify that the values of all the environment variables it needs according to the precedence rules, described in the Base Definitions Volume of IEEE Std 1003.1-8, Chapter 8, Environment Variables, indicate supported locales. If the value of any of these environment-variable searches yields a locale that is not supported (and nonnull), setlocale() shall return a null pointer and the locale of the process shall not be changed. If all environment variables name supported locales, setlocale() shall proceed as if it had been called for each category, using the appropriate value from the associated environment variable or from the implementation-defined default if there is no such value. [/CX]"

Change Number: XSH/TC1/124 [XSH ERN 31,32]

On Page: 1319 Line: 40905-40910 Section: setsockopt

Change From:

"If SO_LINGER is set, the system shall block the process during close() until it can transmit the data or until the time expires. If SO_LINGER is not specified, and close() is issued, the system handles the call in a way that allows the process to continue as quickly as possible."

To:

"If SO_LINGER is set, the system shall block the calling thread during close() until it can transmit the data or until the time expires. If SO_LINGER is not specified, and close() is issued, the system handles the call in a way that allows the calling thread to continue as quickly as possible."

Change Number: XSH/TC1/125 [XSH ERN 129]

On Page: 1330 Line: 41250 Section: shm_open

In the EXAMPLES section

Change From:

"None."

To:

"Creating and mapping a shared memory object

The following code segment demonstrates the use of shm_open() to create a shared memory object which is then sized using ftruncate() before being mapped into the process address space using mmap():

```
#include <unistd.h>
```

```
#include <sys/mman.h>
```

```
...
```

```
#define MAX_LEN 10000
```

```

struct region {          /* Defines "structure" of shared memory */

    int len;

    char buf[MAX_LEN];

};

struct region *rptr;

int fd;

/* Create shared memory object and set its size */

fd = shm_open("/myregion", O_CREAT | O_RDWR, S_IRUSR | S_IWUSR);

if (fd == -1)

    /* Handle error */;

if (ftruncate(fd, sizeof(struct region)) == -1)

    /* Handle error */;

/* Map shared memory object */

rptr = mmap(NULL, sizeof(struct region),

            PROT_READ | PROT_WRITE, MAP_SHARED, fd, 0);

if (rptr == MAP_FAILED)

    /* Handle error */;

/* Now we can refer to mapped region using fields of rptr,

   for example, rptr->len */

..."

```

Change Number: XSH/TC1/126 [XSH ERN 93]

On Page: 1344 Line: 41688-41696 Section: sigaction

In the DESCRIPTION

Change From:

"If the SA_SIGINFO flag (see below) is cleared in the sa_flags field of the sigaction structure, the sa_handler field identifies the action to be associated with the specified signal. [XSI|RTS] If the SA_SIGINFO flag is set in the sa_flags field, and the implementation supports the Realtime Signals Extension option or the XSI Extension option, the sa_sigaction field specifies a signal-catching function. [/XSI|RTS] If the SA_SIGINFO bit is cleared and the sa_handler field specifies a signal-catching function, or if the SA_SIGINFO bit is set, the sa_mask field identifies a set of signals that shall be added to the signal mask of the thread before the signal-catching function is invoked. If the sa_handler field specifies a signal-catching function, the sa_mask field identifies a set of signals that shall be added to the process signal mask before the signal-catching function is invoked."

To:

"If the SA_SIGINFO flag (see below) is cleared in the sa_flags field of the sigaction structure, the sa_handler field identifies the action to be associated with the specified signal. [XSI|RTS] If the SA_SIGINFO flag is set in the sa_flags field, and the implementation supports the Realtime Signals Extension option or the XSI Extension option, the sa_sigaction field specifies a signal-catching function. [/XSI|RTS]

Rationale:

The removed text repeats requirements stated later in the DESCRIPTION section on page 1346 line 41765.

Change Number: XSH/TC1/127 [XSH ERN 33,34,35,36,37]

On Page: 1345-1348 Line: 41736,41761,41843 Section: sigaction

In the DESCRIPTION of SA_SIGINFO, line 41736

Change From:

"The second argument shall point to an object of type siginfo_t explaining the reason why the signal was generated; the third argument can be cast to a pointer to an object of type ucontext_t to refer to the receiving process context that was interrupted when the signal was delivered."

To:

"The second argument shall point to an object of type siginfo_t explaining the reason why the signal was generated; the third argument can be cast to a pointer to an object of type ucontext_t to refer to the receiving thread's context that was interrupted when the signal was delivered."

In the DESCRIPTION of SA_NODEFER, line 41761

Change From:

"If set and sig is caught, sig shall not be added to the process signal mask on entry to the signal handler unless it is included in sa_mask. Otherwise, sig shall always be added to the process signal mask on entry to the signal handler."

To:

"If set and sig is caught, sig shall not be added to the thread's signal mask on entry to the signal handler unless it is included in sa_mask. Otherwise, sig shall always be added to the thread's signal mask on entry to the signal handler."

In the APPLICATION USAGE , line 41843

Change From:

"When the signal handler returns, the receiving process resumes execution at the point it was interrupted unless the signal handler makes other arrangements. If longjmp() or _longjmp() is used to leave the signal handler, then the signal mask must be explicitly restored by the process."

To:

"When the signal handler returns, the receiving thread resumes execution at the point it was interrupted unless the signal handler makes other arrangements. If longjmp() or _longjmp() is used to leave the signal handler, then the signal mask must be explicitly restored."

Change Number: XSH/TC1/128 [XSH ERN 130]

On Page: 1347 Line: 41807 Section: sigaction

In the EXAMPLES section

Change From:

"None."

To:

"Establishing a signal handler

The following example demonstrates the use of sigaction to establish a handler for the SIGINT signal.

```
#include <signal.h>
```

```
static void handler(int signum)

{

    /* Take appropriate actions for signal delivery */

}

int main()

{
```

```
struct sigaction sa;

sa.sa_handler = handler;

sigemptyset(&sa.sa_mask);

sa.sa_flags = SA_RESTART;      /* Restart functions if
                                interrupted by handler */

if (sigaction(SIGINT, &sa, NULL) == -1)

    /* Handle error */;

/* Further code */

}"
```

Change Number: XSH/TC1/129 [XSH ERN 38]

On Page: 1377 Line: 42765-42766 Section: sigtimedwait

Change From:

"An implementation only checks for this error if no signal is pending in set and it is necessary to wait."

To:

"An implementation should only check for this error if no signal is pending in set and it is necessary to wait."

Rationale: This restores the wording from the original base document.

Change Number: XSH/TC1/130 [XSH ERN 94]

On Page: 1380 Line: 42881-42883 Section: sigwait

In the DESCRIPTION section

Change From:

"If more than one thread is using sigwait() to wait for the same signal, no more than one of these threads shall return from sigwait() with the signal number. Which thread returns from sigwait() if more than a single thread is waiting is unspecified."

To:

"If more than one thread is using sigwait() to wait for the same signal, no more than one of these threads shall return from sigwait() with the signal number. If more than a single thread is blocked in sigwait() for a signal when that signal is generated for the process, it is unspecified which of the waiting threads returns from sigwait(). If the signal is generated for a specific thread, as by pthread_kill(), only that thread shall return."

Change Number: XSH/TC1/131 [XSH ERN 38]

On Page: 1388 Line: 43101-43103 Section: sleep

In the RATIONALE section

Change From:

"One is to use the alarm() function to schedule a SIGALRM signal and then suspend the process waiting for that signal."

To:

"One is to use the alarm() function to schedule a SIGALRM signal and then suspend the calling thread waiting for that signal."

Change Number: XSH/TC1/132 [XSH ERN 131]

On Page: 1443 Line: 44720 Section: strptime

In the EXAMPLES section

Change From:

"None."

To:

"Convert a data-plus-time string to broken-down time and then into seconds

The following example demonstrates the use `strptime()` to convert a string into broken-down time. The broken-down time is then converted into seconds since the Epoch using `mktime()`.

```
#include <time.h>
```

```
...
```

```
struct tm tm;
```

```
time_t t;
```

```
if (strptime("6 Dec 2001 12:33:45", "%d %b %Y %H:%M:%S", &tm) == NULL)
```

```
    /* Handle error */;
```

```
printf("year: %d; month: %d; day: %d;\n",
```

```
       tm.tm_year, tm.tm_mon, tm.tm_mday);
```

```
printf("hour: %d; minute: %d; second: %d\n",
```

```
       tm.tm_hour, tm.tm_min, tm.tm_sec);
```

```
printf("week day: %d; year day: %d\n", tm.tm_wday, tm.tm_yday);
```

```
tm.tm_isdst = -1;          /* Not set by strptime(); tells mktime()
```

```
                           to determine if daylight saving time
```

```
                           is in effect */
```

```
t = mktime(&tm);
```

```
if (t == -1)
```

```
    /* Handle error */;
```

```
printf("seconds since the Epoch: %ld\n", (long) t);"
```

Change Number: XSH/TC1/133 [XSH ERN 105]

On Page: 1474 Line: 45617 Section: `sysconf`

In the DESCRIPTION section

Delete the following entry from the table

"_POSIX_SYMLoop_MAX _SC_SYMLoop_MAX"

Rationale:

There is an entry for SYMLoop_MAX in this table, there thus does not need to be an entry for _POSIX_SYMLoop_MAX (the minimum maximum) here.

This corrects an error in the 2003 edition.

Change Number: XSH/TC1/134 [XSH ERN 40,61,62]

On Page: 1474-1476 Line: 45593,45603,45645,45691,45711-45718 Section: `sysconf`

In the DESCRIPTION, line 45593

Delete the following entry from the table

"_POSIX_FILE_LOCKING _SC_FILE_LOCKING"

In the DESCRIPTION, line 45603

Delete the following entry from the table

"_POSIX_MULTI_PROCESS _SC_MULTI_PROCESS"

In the DESCRIPTION, line 45645

Delete the following entry from the table

"_POSIX2_C_VERSION _SC_2_C_VERSION"

In the DESCRIPTION, line 45691

Delete the following entry from the table

"_XOPEN_XCU_VERSION _SC_XOPEN_XCU_VERSION"

In the APPLICATION USAGE, lines 45711-45718

Delete the paragraph

"If the value of `sysconf(_SC_2_VERSION)` is not equal to the value of the `_POSIX2_VERSION` symbolic constant, the utilities available via `system()` or `popen()` might not behave as described in the Shell and Utilities volume of IEEE Std 1003.1-2001. This would mean that the application is not running in an environment that conforms to the Shell and Utilities volume of IEEE Std 1003.1-2001. Some applications might be able to deal with this, others might not. However, the functions defined in this volume of IEEE Std 1003.1-2001 continue to operate as specified, even if `sysconf(_SC_2_VERSION)` reports that the utilities no longer perform as specified."

Rationale: `_POSIX_FILE_LOCKING` and `_POSIX_MULTI_PROCESS` were erroneous entries. `_POSIX2_C_VERSION` was removed by P1003.1a/P1003.2b and should have been removed here too. `_XOPEN_XCU_VERSION` should have been removed as per The Open Group Base Working Group resolution bwg2001-006.

Change Number: XSH/TC1/135 [XBD ERN 27]

On Page: 1474 Line: 45616,45632,45635,45701,45718 Section: `sysconf`

In the DESCRIPTION section:

Add into the table after the `_POSIX_SPORADIC_SERVER` entry:

"_POSIX_SS_REPL_MAX _SC_SS_REPL_MAX"

Add into the table after the `_POSIX_TRACE_EVENT_FILTER` entry:

"_POSIX_TRACE_EVENT_NAME_MAX _SC_TRACE_EVENT_NAME_MAX"

Add into the table after the `_POSIX_TRACE_MEMORY_OBJECTS` entry:

"_POSIX_TRACE_NAME_MAX _SC_TRACE_NAME_MAX"

"_POSIX_TRACE_SYS_MAX _SC_TRACE_SYS_MAX"

"_POSIX_TRACE_USER_EVENT_MAX _SC_TRACE_USER_EVENT_MAX"

Add to the end of the RETURN VALUE section as a new paragraph:

"If the variable corresponding to name is dependent on an unsupported option the results are unspecified."

Add to the end of the APPLICATION USAGE section:

"Application writers should check whether an option, such as `_POSIX_TRACE` is supported, prior to obtaining and using values for related variables such as `_POSIX_TRACE_NAME_MAX`."

Change Number: XSH/TC1/136 [XSH ERN 107]

On Page: 1475 Line: 45659 Section: `sysconf`

In the DESCRIPTION section

Delete the following entry from the table:

"_REGEX_VERSION _SC_REGEX_VERSION"

Rationale:

The text in P1003.1a indicates that this was to be used for versioning information related to the `_SC_REGEX_VERSION` (sic)

(.1a D17): "172 POSIX.2RE: The version of ISO/IEC 9945-2 that defines `_POSIX2_VERSION` to be the value indicated by the return value of `sysconf(_SC_REGEX_VERSION)`."

Given that the approach for `"_SC_2_C_VERSION"` and `"_SC_XOPEN_XCU_VERSION"` has been to remove them since they are unnecessary this is removed for consistency.

Change Number: XSH/TC1/137 [XSH ERN 95]

On Page: 1521,1522 Line: 47034,47050 Section: `timer_create`

In the DESCRIPTION section

Add as a new last paragraph, shaded and margin marked "TSA":

"If `evp->sigev_sigev_notify` is `SIGEV_THREAD` and `sev->sigev_notify_attributes` is not NULL, if the attribute pointed to by `sev->sigev_notify_attributes` has a thread stack address specified by a call to `pthread_attr_setstack()` or `pthread_attr_setstackaddr()` the results are unspecified if the signal is generated more than once."

In the APPLICATION USAGE section:

Change From:

"None."

To:

"If a timer is created which has `evp->sigev_sigev_notify` set to `SIGEV_THREAD` and the attribute pointed to by `evp->sigev_notify_attributes` has a thread stack address specified by a call to either `pthread_attr_setstack()` or `pthread_attr_setstackaddr()`, the memory dedicated as a thread stack can not be recovered. The reason for this is that the threads created in response to a timer expiration are created detached, or in an unspecified way if the thread attribute's `detachstate` is `PTHREAD_CREATE_JOINABLE`. In neither case is it valid to call `pthread_join()`, which makes it impossible to determine the lifetime of the created thread which thus means the stack memory can not be reused."

Change Number: XSH/TC1/138 [XSH ERN 74]

On Page: 1524 Line: 47131 Section: `timer_delete`

In the ERRORS section

Change From:

"The `timer_delete()` function shall fail if:"

To:

"The `timer_delete()` function may fail if:"

Change Number: XSH/TC1/139 [XSH ERN 75]

On Page: 1526 Line: 47211-47213 Section: `timer_getoverrun`

Change From:

"The `timer_getoverrun()`, `timer_gettime()`, and `timer_settime()` functions shall fail if:

[EINVAL] The `timerid` argument does not correspond to an ID returned by `timer_create()` but not yet deleted by `timer_delete()`."

To:

"The `timer_getoverrun()`, `timer_gettime()`, and `timer_settime()` functions may fail if:

[EINVAL] The `timerid` argument does not correspond to an ID returned by `timer_create()` but not yet deleted by `timer_delete()`."

Notes to Editors:

Move the revised text section after the current following [EINVAL] error so that the "shall fail" case occurs before the "may fail" cases.

Change Number: XSH/TC1/140 [XSH ERN 95]

On Page: 1526 Line: 47217,47221 Section: `timer_getoverrun`

In the ERRORS section

Add the following additional text

"The `timer_settime()` function may fail if:

[EINVAL] The `it_interval` member of value is not zero and the timer was created with notification by creation of a new thread (`sigev_sigev_notify` was `SIGEV_THREAD`) and a fixed stack address has been set in the thread attribute pointed to by `sigev_notify_attributes`."

In the APPLICATION USAGE section

Change From:

"None."

To:

"Using fixed stack addresses is problematic when timer expiration is signalled by the creation of a new thread. Since it can not be assumed that the thread created for one expiration is finished before the next expiration of the timer, it could happen that two threads use the same memory as stack at the same time. This is invalid and produces undefined results."

Change Number: XSH/TC1/141 [XSH ERN 137]

On Page: 1533 Line: 47433-47434 Section: `tempnam`

In the DESCRIPTION section

Change From:

"The implementation shall behave as if no function defined in this volume of IEEE Std 1003.1-2001 calls `tmpnam()`."

To:

"The implementation shall behave as if no function defined in this volume of IEEE Std 1003.1-2001 except `tempnam()` calls `tmpnam()`."

Change Number: XSH/TC1/142 [XSH ERN 52,98]

On Page: 1555 Line: 48027-48028,48051 Section: umask

In the DESCRIPTION section

Change From:

"The process' file mode creation mask is used during open(), creat(), mkdir(), and mkfifo() to turn off permission bits in the mode argument supplied."

To:

"The process' file mode creation mask is used during open(), creat(), mkdir(), mkfifo(), mknod(), mq_open() and sem_open() to turn off permission bits in the mode argument supplied."

(Ed note:mq_open is shaded and margin marked MSG, sem_open is shaded and margin marked SEM, mknod is shaded and margin marked XSI)

In the SEE ALSO section

Add "mknod()", mq_open()" and "sem_open()" in alphabetical order.

Change Number: XSH/TC1/143 [XSH ERN 41]

On Page: 1568 Line: 48450-48452,48465 Section: usleep

Change From:

"It is also unspecified whether the SIGALRM signal is blocked, unless the process signal mask is restored as part of the environment."

To:

"It is also unspecified whether the SIGALRM signal is blocked, unless the thread's signal mask is restored as part of the environment."

Add as a new final paragraph to the DESCRIPTION:

"The usleep() function need not be reentrant. A function that is not required to be reentrant is not required to be thread-safe."

Change Number: XSH/TC1/144 [XSH ERN 116]

On Page: 1592 Line: 49061 Section: wait

In the EXAMPLES section

Change From:

"None."

To:

"Waiting for a child process and then checking its status

The following example demonstrates the use of waitpid(), fork(), and the macros used to interpret the status value returned by waitpid() (and wait()). The code segment creates a child process which does some unspecified work. Meanwhile the parent loops performing calls to waitpid() to monitor the status of the child. The loop terminates when child termination is detected.

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
#include <unistd.h>
```

```
#include <sys/wait.h>
```

```
...
```

```
pid_t child_pid, wpid;
```

```
int status;
```

```
child_pid = fork();
```

```
if (child_pid == -1) {          /* fork() failed */
```

```
    perror("fork");
```

```
    exit(EXIT_FAILURE);
```

```

}

if (child_pid == 0) {          /* This is the child */

    /* Child does some work and then terminates */

    ...

} else {                      /* This is the parent */

    do {

        wpid = waitpid(child_pid, &status, WUNTRACED

#ifdef WCONTINUED             /* Not all implementations support this */

                                | WCONTINUED

#endif

                                );

        if (wpid == -1) {

            perror("waitpid");

            exit(EXIT_FAILURE);

        }

        if (WIFEXITED(status)) {

            printf("child exited, status=%d\n", WEXITSTATUS(status));

        } else if (WIFSIGNALED(status)) {

            printf("child killed (signal %d)\n", WTERMSIG(status));

        } else if (WIFSTOPPED(status)) {

            printf("child stopped (signal %d)\n", WSTOPSIG(status));

#ifdef WIFCONTINUED           /* Not all implementations support this */

        } else if (WIFCONTINUED(status)) {

            printf("child continued\n");

#endif

    }

    } else {                  /* Non-standard case -- may never happen */

        printf("Unexpected status (0x%x)\n", status);

    }

    } while (!WIFEXITED(status) && !WIFSIGNALED(status));

}

```

Change Number: XSH/TC1/145 [XSH ERN 166]

On Page: 1659 Line: 51172-51173 Section: write

In the ERRORS section

Change From:

"In the latter case, if the socket is of type SOCK_STREAM, the SIGPIPE signal is generated to the calling process."

To:

"In the latter case, if the socket is of type SOCK_STREAM, a SIGPIPE signal shall also be sent to the thread."

Change Number: XSH/TC1/146 [XSH ERN 42]

On Page: 1660 Line: 51223 Section: write

In the RATIONALE section

Change From:

"Otherwise, the process may block; that is, pause until enough space is available for writing."

To:

"Otherwise, the calling thread may block; that is, pause until enough space is available for writing."

Change Number: XSH/TC1/147 [XSH ERN 108]

On Page: 1667 Line: 51442,51453 Section: y0

In the RETURN VALUE section

Change From:

"If x is 0.0, -HUGE_VAL shall be returned and a range error may occur".

To:

"If x is 0.0, -HUGE_VAL shall be returned and a pole error may occur."

In the ERRORS section

Insert before line 51453

"Pole error The value of x is zero.

If the integer expression (math_errhandling & MATH_ERRNO) is non-zero, then errno shall be set to [ERANGE]. If the integer expression (math_errhandling & MATH_ERREXCEPT) is non-zero, then the divide-by-zero floating-point exception shall be raised."

In the ERRORS section, line 51453

Change From:

"Range Error The value of x is 0.0, or the correct result would cause overflow."

To:

"Range Error The correct result would cause overflow."

Rationale:

The changes are made for consistency with the general rules stated in "Treatment of Error Conditions for Mathematical Functions" in XBD.