# INTERNATIONAL STANDARD

## **ISO/IEC** 8652

Third edition 2012-12-15

# Information technology — Programming languages — Ada

Technologies de l'information — Langages de programmation — Ada



Reference number ISO/IEC 8652:2012(E)



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Published in Switzerland

### **Table of Contents**

Table of Contents	i
Foreword	X
Introduction	χi
1 General	1
1.1 Scope	. 1
1.1.1 Extent	
1.1.2 Structure	. 2
1.1.3 Conformity of an Implementation with the Standard	. 4
1.1.4 Method of Description and Syntax Notation	. 5
1.1.5 Classification of Errors	. 6
1.2 Normative References	
1.3 Terms and Definitions	. 8
2 Lexical Elements	9
2.1 Character Set	
2.2 Lexical Elements, Separators, and Delimiters	
2.3 Identifiers	
2.4 Numeric Literals	13
2.4.1 Decimal Literals	
2.4.2 Based Literals	13
2.5 Character Literals	14
2.6 String Literals	14
2.7 Comments	15
2.8 Pragmas	15
2.9 Reserved Words	17
3 Declarations and Types	
3.1 Declarations	19
3.2 Types and Subtypes	20
3.2.1 Type Declarations	
3.2.2 Subtype Declarations	
3.2.3 Classification of Operations	
3.2.4 Subtype Predicates	24
3.3 Objects and Named Numbers	26
3.3.1 Object Declarations	28
3.3.2 Number Declarations	31
3.4 Derived Types and Classes	31
3.4.1 Derivation Classes	34 25
3.5 Scalar Types	30 30
3.5.2 Character Types	40 11
3.5.3 Boolean Types	
3.5.4 Integer Types	-74 42
3.5.5 Operations of Discrete Types	
3.5.6 Real Types	4F
3.5.7 Floating Point Types	46
3.5.8 Operations of Floating Point Types	
3.5.9 Fixed Point Types	
3.5.10 Operations of Fixed Point Types	

3.6 Array Types	51
3.6.1 Index Constraints and Discrete Ranges	
3.6.2 Operations of Array Types	
3.6.3 String Types	
3.7 Discriminants	
3.7.1 Discriminant Constraints	
3.7.2 Operations of Discriminated Types	
3.8 Record Types	
3.8.1 Variant Parts and Discrete Choices	
3.9 Tagged Types and Type Extensions	
3.9.1 Type Extensions	
3.9.2 Dispatching Operations of Tagged Types	
3.9.3 Abstract Types and Subprograms	
3.9.4 Interface Types	
3.10 Access Types	
3.10.1 Incomplete Type Declarations	
3.10.2 Operations of Access Types	79
3.11 Declarative Parts	85
3.11.1 Completions of Declarations	85
4 Names and Expressions	87
4.1 Names	87
4.1.1 Indexed Components	
4.1.2 Slices	
4.1.3 Selected Components	89
4.1.4 Attributes	91
4.1.5 User-Defined References	92
4.1.6 User-Defined Indexing	93
4.2 Literals	95
4.3 Aggregates	96
4.3.1 Record Aggregates	96
4.3.2 Extension Aggregates	98
4.3.3 Array Aggregates	99
4.3.3 Array Aggregates4.4 Expressions	102
4.5 Operators and Expression Evaluation	103
4.5.1 Logical Operators and Short-circuit Contr	rol Forms104
4.5.2 Relational Operators and Membership Te	
4.5.3 Binary Adding Operators	109
4 5 4 Unary Adding Operators	110
4.5.4 Unary Adding Operators 4.5.5 Multiplying Operators	110
4.5.6 Highest Precedence Operators	112
4.5.7 Conditional Expressions	112
4.5.8 Quantified Expressions	
4.6 Type Conversions	
4.7 Qualified Expressions	
4.8 Allocators	
4.9 Static Expressions and Static Subtypes	122
4.9.1 Statically Matching Constraints and Subt	
5 Statements	
5.1 Simple and Compound Statements - Sequence	ces of Statements 197
5.2 Assignment Statements	
5.3 If Statements	
5.4 Case Statements	
v. 7 vast statellitilis	

	5.5 Loop Statements	13	3
	5.5.2 Generalized Loop Iteration		
	5.6 Block Statements		
	5.7 Exit Statements		
	5.8 Goto Statements	13	7
6			
	6.1 Subprogram Declarations		
	6.1.1 Preconditions and Postconditions		
	6.2 Formal Parameter Modes		
	6.3 Subprogram Bodies		
	6.3.1 Conformance Rules		
	6.3.2 Inline Expansion of Subprograms		
	6.4 Subprogram Calls		
	6.4.1 Parameter Associations		
	6.5 Return Statements		
	6.5.1 Nonreturning Procedures		
	6.6 Overloading of Operators		
	6.7 Null Procedures		
	6.8 Expression Functions		
7	Packages	15	9
	7.1 Package Specifications and Declarations	15	9
	7.2 Package Bodies		
	7.3 Private Types and Private Extensions	16	1
	7.3.1 Private Operations	16	3
	7.3.2 Type Invariants	16	5
	7.4 Deferred Constants	16	7
	7.5 Limited Types	16	8
	7.6 Assignment and Finalization	17	0
	7.6.1 Completion and Finalization	17	2
8	Visibility Rules	17	5
	8.1 Declarative Region	17	5
	8.2 Scope of Declarations		
	8.3 Visibility	17	7
	8.3.1 Overriding Indicators		
	8.4 Use Clauses	18	0
	8.5 Renaming Declarations	18	1
	8.5.1 Object Renaming Declarations	18	2
	8.5.2 Exception Renaming Declarations	18	3
	8.5.3 Package Renaming Declarations		
	8.5.4 Subprogram Renaming Declarations	18	4
	8.5.5 Generic Renaming Declarations	18	6
	8.6 The Context of Overload Resolution	18	6
9	Tasks and Synchronization		
	9.1 Task Units and Task Objects	18	9
	9.2 Task Execution - Task Activation	19	2
	9.3 Task Dependence - Termination of Tasks	19	3
	9.4 Protected Units and Protected Objects	19	4
	9.5 Intertask Communication		
	9.5.1 Protected Subprograms and Protected Actions		
	9.5.2 Entries and Accent Statements	20	n

9.5.3 Entry Calls	.203
9.5.4 Requeue Statements	.205
9.6 Delay Statements, Duration, and Time	.207
9.6.1 Formatting, Time Zones, and other operations for Time	.209
9.7 Select Statements	
9.7.1 Selective Accept	
9.7.2 Timed Entry Calls	
9.7.3 Conditional Entry Calls	218
9.7.4 Asynchronous Transfer of Control	.219
9.8 Abort of a Task - Abort of a Sequence of Statements	
9.9 Task and Entry Attributes	
9.10 Shared Variables	
9.11 Example of Tasking and Synchronization	
10 Program Structure and Compilation Issues	225
10.1 Separate Compilation	.225
10.1.1 Compilation Units - Library Units	
10.1.2 Context Clauses - With Clauses	.228
10.1.3 Subunits of Compilation Units	
10.1.4 The Compilation Process	
10.1.5 Pragmas and Program Units	
10.1.6 Environment-Level Visibility Rules	
10.2 Program Execution	
10.2.1 Elaboration Control	.236
11 Exceptions	241
11.1 Exception Declarations	244
11.2 Exception Handlers	241
11.3 Raise Statements	
11.4 Exception Handling	
11.4.1 The Package Exceptions	243
11.4.2 Pragmas Assert and Assertion_Policy	244
11.4.3 Example of Exception Handling	
11.5 Suppressing Checks	249
11.6 Exceptions and Optimization	252
12 Generic Units	
12.1 Generic Declarations	
12.2 Generic Bodies	
12.3 Generic Instantiation	.256
12.4 Formal Objects	.258
12.5 Formal Types	.259
12.5.1 Formal Private and Derived Types	.261
12.5.2 Formal Scalar Types	.263
12.5.3 Formal Array Types12.5.4 Formal Access Types	.263
12.5.5 Formal Interface Types	.204
12.6 Formal Subprograms	
12.7 Formal Packages	
12.8 Example of a Generic Package	
·	
13 Representation Issues	
13.1 Operational and Representation Aspects	
13.1.1 Aspect Specifications	
13.2 Packed Types	.278

13.3 Operational and Representation Attributes	279
13.4 Enumeration Representation Clauses	285
13.5 Record Layout	
13.5.1 Record Representation Clauses	
13.5.2 Storage Place Attributes	
13.5.3 Bit Ordering	
13.6 Change of Representation	
13.7 The Package System	
13.7.1 The Package System.Storage_Elements	
13.7.2 The Package System.Address_To_Access_Conversions	
13.8 Machine Code Insertions	
13.9 Unchecked Type Conversions	
13.9.1 Data Validity	
13.9.2 The Valid Attribute	
13.10 Unchecked Access Value Creation	
13.11 Storage Management	
13.11.1 Storage Allocation Attributes	
13.11.2 Unchecked Storage Deallocation	
13.11.3 Default Storage Pools	
13.11.4 Storage Subpools	304
13.11.5 Subpool Reclamation	306
13.11.6 Storage Subpool Example	307
13.12 Pragma Restrictions and Pragma Profile	309
13.12.1 Language-Defined Restrictions and Profiles	
13.13 Streams	
13.13.1 The Package Streams	
13.13.2 Stream-Oriented Attributes	
13.14 Freezing Rules	
he Standard Libraries	
ne Standard Libraries	3 <b>Z</b> I
nnex A (normative) Predefined Language Environment	323
A.1 The Package Standard	
A.2 The Package Ada	
A.3 Character Handling	
A.3.1 The Packages Characters, Wide_Characters, and Wide_Wide_Characters	
A.3.2 The Package Characters.HandlingA.3.3 The Package Characters.Latin_1	331
A.3.3 The Package Characters.Latin_1	333
A.3.4 The Package Characters.Conversions	338
A.3.5 The Package Wide_Characters.Handling	340
A.3.6 The Package Wide_Wide_Characters.Handling	342
A.4 String HandlingA.4.1 The Package Strings	343
A.4.1 The Package Strings	343
A.4.2 The Package Strings.Maps	343
A.4.3 Fixed-Length String Handling	
A.4.4 Bounded-Length String Handling	
A.4.5 Unbounded-Length String Handling	361
A.4.6 String-Handling Sets and Mappings	200
A.4.7 Wide_String Handling	300
A.4.8 Wide_Wide_String Handling	366
A.4.8 Wide_Wide_String Handling A.4.9 String Hashing	366 368
	366 368 371
A.4.9 String Hashing	366 368 371 372

A.5.1 Elementary Functions	378
A.5.2 Random Number Generation	381
A.5.3 Attributes of Floating Point Types	386
A.5.4 Attributes of Fixed Point Types	
A.6 Input-Output	390
A.7 External Files and File Objects	
A.8 Sequential and Direct Files	
A.8.1 The Generic Package Sequential_IO	
A.8.2 File Management	
A.8.3 Sequential Input-Output Operations	
A.8.4 The Generic Package Direct_IO	
A.8.5 Direct Input-Output Operations	
A.9 The Generic Package Storage_IO	
A.10 Text Input-Output	
A.10.1 The Package Text IO	
A.10.2 Text File Management	
A.10.3 Default Input, Output, and Error Files	
A.10.4 Specification of Line and Page Lengths	
A.10.5 Operations on Columns, Lines, and Pages	
A.10.6 Get and Put Procedures	
A.10.7 Input-Output of Characters and Strings	
A.10.8 Input-Output for Integer Types	
A.10.9 Input-Output for Real Types	
A.10.10 Input-Output for Enumeration Types	
A.10.11 Input-Output for Bounded Strings	
A.10.12 Input-Output for Unbounded Strings	
A.11 Wide Text Input-Output and Wide Wide Text Input-Output	
A.12 Stream Input-Output	
A.12.1 The Package Streams.Stream_IO	
A.12.2 The Package Text_IO.Text_Streams	
A.12.3 The Package Wide_Text_IO.Text_Streams	
A.12.4 The Package Wide_Wide_Text_IO.Text_Streams	
A.13 Exceptions in Input-Output	423
A.14 File Sharing	425
A.15 The Package Command_Line  A.16 The Package Directories	425
A.16 The Package Directories	426
A.16.1 The Package Directories.Hierarchical_File_Names	433
A.17 The Package Environment_Variables	435
A.18 Containers	438
A.18.1 The Package Containers	438
A.18.1 The Package ContainersA.18.2 The Generic Package Containers.Vectors	438
A.18.3 The Generic Package Containers.Doubly_Linked_Lists	454
A.18.4 Maps	465
A.18.5 The Generic Package Containers.Hashed_Maps	471
A.18.6 The Generic Package Containers.Ordered_Maps	475
A.18.7 Sets	
A.18.8 The Generic Package Containers.Hashed_Sets	
A.18.9 The Generic Package Containers.Ordered Sets	
A.18.10 The Generic Package Containers.Multiway_TreesTrees	
A.18.11 The Generic Package Containers.Indefinite_Vectors	
A.18.12 The Generic Package Containers.Indefinite Doubly Linked Lists	
A.18.13 The Generic Package Containers.Indefinite_Hashed_Maps	
A.18.14 The Generic Package Containers.Indefinite_Ordered_Maps	

A.18.16 The Generic Package Containers.Indefinite_Ordered_Sets	512
A.18.17 The Generic Package Containers.Indefinite_Multiway_Trees	
A.18.18 The Generic Package Containers.Indefinite_Holders	
A.18.19 The Generic Package Containers.Bounded_Vectors	
A.18.20 The Generic Package Containers.Bounded_Doubly_Linked_Lists	
A.18.21 The Generic Package Containers.Bounded_Hashed_Maps	
A.18.22 The Generic Package Containers.Bounded_Ordered_Maps	
A.18.23 The Generic Package Containers.Bounded_Hashed_Sets	
A.18.24 The Generic Package Containers.Bounded Ordered Sets	
A.18.25 The Generic Package Containers.Bounded_Multiway_Trees	
	524
A.18.27 The Generic Package Containers.Synchronized_Queue_Interfaces	
A.18.28 The Generic Package Containers. Unbounded Synchronized Queues	
A.18.29 The Generic Package Containers.Bounded_Synchronized_Queues	
A.18.30 The Generic Package Containers. Unbounded_Priority_Queues	
A.18.31 The Generic Package Containers.Bounded_Priority_Queues	
A.18.32 Example of Container Use	
A.19 The Package Locales	
Annex B (normative) Interface to Other Languages	533
B.1 Interfacing Aspects	533
B.2 The Package Interfaces	536
B.3 Interfacing with C and C++	
B.3.1 The Package Interfaces.C.Strings	
B.3.2 The Generic Package Interfaces.C.Pointers	
B.3.3 Unchecked Union Types	
B.4 Interfacing with COBOL	
B.5 Interfacing with Fortran	
Annex C (normative) Systems Programming	
C.1 Access to Machine Operations	
C.2 Required Representation Support	
C.3 Interrupt Support	
C.3.1 Protected Procedure Handlers	
C.3.2 The Package Interrupts	
	566
C.5 Pragma Discard_Names C.6 Shared Variable Control	566
C.6 Shared Variable Control	567
C.7 Task Information	569
C.7 Task Information C.7.1 The Package Task_Identification	569
C.7.2 The Package Task_Attributes C.7.3 The Package Task_Termination	571
C.7.3 The Package Task_Termination	573
Annex D (normative) Real-Time Systems	575
D.1 Task Priorities	
D.2 Priority Scheduling	577
D.2.1 The Task Dispatching Model	577
D.2.2 Task Dispatching Pragmas	
D.2.3 Preemptive Dispatching	580
D.2.4 Non-Preemptive Dispatching	580
D.2.5 Round Robin Dispatching	
D.2.6 Earliest Deadline First Dispatching	
D.3 Priority Ceiling Locking	
D.4 Entry Queuing Policies	
D.5 Dynamic Priorities	
L L. W. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1.	-

D.5.1 Dynamic Priorities for Tasks	
D.5.2 Dynamic Priorities for Protected Objects	
D.6 Preemptive Abort	
D.7 Tasking Restrictions	
D.8 Monotonic Time	
D.9 Delay Accuracy	596
D.10 Synchronous Task Control	597
D.10.1 Synchronous Barriers	598
D.11 Asynchronous Task Control	599
D.12 Other Optimizations and Determinism Rules	600
D.13 The Ravenscar Profile	
D.14 Execution Time	602
D.14.1 Execution Time Timers	604
D.14.2 Group Execution Time Budgets	606
D.14.3 Execution Time of Interrupt Handlers	
D.15 Timing Events	
D.16 Multiprocessor Implementation	
D.16.1 Multiprocessor Dispatching Domains	
Annex E (normative) Distributed Systems	
E.1 Partitions	
E.2 Categorization of Library Units	
E.2.1 Shared Passive Library Units	617
E.2.2 Remote Types Library Units	618
E.2.3 Remote Call Interface Library Units	619
E.3 Consistency of a Distributed System	
E.4 Remote Subprogram Calls	621
E.4.1 Asynchronous Remote Calls	623
E.4.2 Example of Use of a Remote Access-to-Class-Wide Type	623
E.5 Partition Communication Subsystem	625
Annex F (normative) Information Systems	629
F.1 Machine_Radix Attribute Definition Clause	629
F.2 The Package Decimal	629
F.3 Edited Output for Decimal Types	630
F.3.1 Picture String Formation	
F.3.2 Edited Output Generation	635
F.3.2 Edited Output GenerationF.3.3 The Package Text_IO.Editing	638
F.3.4 The Package Wide_Text_IO.Editing	641
F.3.5 The Package Wide_Wide_Text_IO.Editing	642
Annex G (normative) Numerics	642
G.1 Complex Arithmetic	643
G.1.1 Complex Types	
G.1.2 Complex Elementary Functions	647
G.1.3 Complex Input-Output	651
G.1.4 The Package Wide_Text_IO.Complex_IO	
G.1.5 The Package Wide_Wide_Text_IO.Complex_IO	
G.2 Numeric Performance Requirements	
G.2.1 Model of Floating Point Arithmetic	
G.2.2 Model-Oriented Attributes of Floating Point Types	
G.2.3 Model of Fixed Point Arithmetic	
G.2.4 Accuracy Requirements for the Elementary Functions	
G.2.5 Performance Requirements for Random Number Generation	659
G.2.6 Accuracy Requirements for Complex Arithmetic	660

G.3 Vector and Matrix Manipulation	662
G.3.1 Real Vectors and Matrices	662
G.3.2 Complex Vectors and Matrices	667
Annex H (normative) High Integrity Systems	
H.1 Pragma Normalize_Scalars	
H.2 Documentation of Implementation Decisions	
H.3 Reviewable Object Code	
H.3.1 Pragma Reviewable	
H.3.2 Pragma Inspection_Point	
H.4 High Integrity Restrictions	
H.5 Pragma Detect_Blocking	
H.6 Pragma Partition_Elaboration_Policy	
Annex J (normative) Obsolescent Features	
J.1 Renamings of Library Units	
J.2 Allowed Replacements of Characters	
J.3 Reduced Accuracy Subtypes	
J.4 The Constrained Attribute	
J.5 ASCII	
J.6 Numeric_Error	
J.7 At Clauses	687
J.7.1 Interrupt Entries	688
J.8 Mod Clauses	689
J.9 The Storage_Size Attribute	689
J.10 Specific Suppression of Checks	689
J.11 The Class Attribute of Untagged Incomplete Types	690
J.12 Pragma Interface	690
J.13 Dependence Restriction Identifiers	690
J.14 Character and Wide_Character Conversion Functions	
J.15 Aspect-related Pragmas	
J.15.1 Pragma Inline	691
J.15.2 Pragma No_Return	
J.15.3 Pragma Pack	
J.15.4 Pragma Storage Size	
J.15.5 Interfacing Pragmas	
J.15.6 Pragma Unchecked_Union	693
J.15.7 Pragmas Interrupt_Handler and Attach_Handler	
.I 15 8 Shared Variable Pragmas	694
J.15.8 Shared Variable Pragmas J.15.9 Pragma CPU	69 <i>F</i>
J 15 10 Pragma Dispatching Domain	695
J 15 11 Pragmas Priority and Interrunt Priority	696
J.15.10 Pragma Dispatching_Domain	696
J.15.13 Pragma Asynchronous	697
Annex K (informative) Language-Defined Aspects and Attributes	
K.1 Language-Defined Aspects	
K.1 Language-Defined AspectsK.2 Language-Defined Attributes	. 699 702
Annex L (informative) Language-Defined Pragmas	717
Annex M (informative) Summary of Documentation Requirements	719
M.1 Specific Documentation Requirements	719
M.2 Implementation-Defined Characteristics	
M.3 Implementation Advice	726

Annex N (informative) Glossary	
Annex P (informative) Syntax Summary	741
Annex Q (informative) Language-Defined Entities	
Q.1 Language-Defined Packages	769
Q.2 Language-Defined Types and Subtypes	771
Q.3 Language-Defined Subprograms	
Q.4 Language-Defined ExceptionsQ.5 Language-Defined Objects	
Index	
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### **Foreword**

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.

International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 2.

The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

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International Standard ISO/IEC 8652 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information Technology* Subcommittee SC22, *Programming languages, their environments and system software interfaces*.

This third edition cancels and replaces the second edition (ISO/IEC 8652:1995), which has been technically revised. It also incorporates the Technical Corrigendum ISO/IEC 8652:1995:COR.1:2001 and Amendment ISO/IEC 8652:1995:AMD 1:2007.

### Introduction

#### **Design Goals**

Ada was originally designed with three overriding concerns: program reliability and maintenance, programming as a human activity, and efficiency. The 1995 revision to the language was designed to provide greater flexibility and extensibility, additional control over storage management and synchronization, and standardized packages oriented toward supporting important application areas, while at the same time retaining the original emphasis on reliability, maintainability, and efficiency. This third edition provides further flexibility and adds more standardized packages within the framework provided by the 1995 revision.

The need for languages that promote reliability and simplify maintenance is well established. Hence emphasis was placed on program readability over ease of writing. For example, the rules of the language require that program variables be explicitly declared and that their type be specified. Since the type of a variable is invariant, compilers can ensure that operations on variables are compatible with the properties intended for objects of the type. Furthermore, error-prone notations have been avoided, and the syntax of the language avoids the use of encoded forms in favor of more English-like constructs. Finally, the language offers support for separate compilation of program units in a way that facilitates program development and maintenance, and which provides the same degree of checking between units as within a unit.

Concern for the human programmer was also stressed during the design. Above all, an attempt was made to keep to a relatively small number of underlying concepts integrated in a consistent and systematic way while continuing to avoid the pitfalls of excessive involution. The design especially aims to provide language constructs that correspond intuitively to the normal expectations of users.

Like many other human activities, the development of programs is becoming ever more decentralized and distributed. Consequently, the ability to assemble a program from independently produced software components continues to be a central idea in the design. The concepts of packages, of private types, and of generic units are directly related to this idea, which has ramifications in many other aspects of the language. An allied concern is the maintenance of programs to match changing requirements; type extension and the hierarchical library enable a program to be modified while minimizing disturbance to existing tested and trusted components.

No language can avoid the problem of efficiency. Languages that require over-elaborate compilers, or that lead to the inefficient use of storage or execution time, force these inefficiencies on all machines and on all programs. Every construct of the language was examined in the light of present implementation techniques. Any proposed construct whose implementation was unclear or that required excessive machine resources was rejected.

#### **Language Summary**

An Ada program is composed of one or more program units. Program units may be subprograms (which define executable algorithms), packages (which define collections of entities), task units (which define concurrent computations), protected units (which define operations for the coordinated sharing of data between tasks), or generic units (which define parameterized forms of packages and subprograms). Each program unit normally consists of two parts: a specification, containing the information that must be visible to other units, and a body, containing the implementation details, which need not be visible to other units. Most program units can be compiled separately.

This distinction of the specification and body, and the ability to compile units separately, allows a program to be designed, written, and tested as a set of largely independent software components.

An Ada program will normally make use of a library of program units of general utility. The language provides means whereby individual organizations can construct their own libraries. All libraries are structured in a hierarchical manner; this enables the logical decomposition of a subsystem into

individual components. The text of a separately compiled program unit must name the library units it requires.

#### Program Units

A subprogram is the basic unit for expressing an algorithm. There are two kinds of subprograms: procedures and functions. A procedure is the means of invoking a series of actions. For example, it may read data, update variables, or produce some output. It may have parameters, to provide a controlled means of passing information between the procedure and the point of call. A function is the means of invoking the computation of a value. It is similar to a procedure, but in addition will return a result.

A package is the basic unit for defining a collection of logically related entities. For example, a package can be used to define a set of type declarations and associated operations. Portions of a package can be hidden from the user, thus allowing access only to the logical properties expressed by the package specification.

Subprogram and package units may be compiled separately and arranged in hierarchies of parent and child units giving fine control over visibility of the logical properties and their detailed implementation.

A task unit is the basic unit for defining a task whose sequence of actions may be executed concurrently with those of other tasks. Such tasks may be implemented on multicomputers, multiprocessors, or with interleaved execution on a single processor. A task unit may define either a single executing task or a task type permitting the creation of any number of similar tasks.

A protected unit is the basic unit for defining protected operations for the coordinated use of data shared between tasks. Simple mutual exclusion is provided automatically, and more elaborate sharing protocols can be defined. A protected operation can either be a subprogram or an entry. A protected entry specifies a Boolean expression (an entry barrier) that must be True before the body of the entry is executed. A protected unit may define a single protected object or a protected type permitting the creation of several similar objects.

#### Declarations and Statements

The body of a program unit generally contains two parts: a declarative part, which defines the logical entities to be used in the program unit, and a sequence of statements, which defines the execution of the program unit.

The declarative part associates names with declared entities. For example, a name may denote a type, a constant, a variable, or an exception. A declarative part also introduces the names and parameters of other nested subprograms, packages, task units, protected units, and generic units to be used in the program unit.

The sequence of statements describes a sequence of actions that are to be performed. The statements are executed in succession (unless a transfer of control causes execution to continue from another place).

An assignment statement changes the value of a variable. A procedure call invokes execution of a procedure after associating any actual parameters provided at the call with the corresponding formal parameters.

Case statements and if statements allow the selection of an enclosed sequence of statements based on the value of an expression or on the value of a condition.

The loop statement provides the basic iterative mechanism in the language. A loop statement specifies that a sequence of statements is to be executed repeatedly as directed by an iteration scheme, or until an exit statement is encountered.

A block statement comprises a sequence of statements preceded by the declaration of local entities used by the statements.

Certain statements are associated with concurrent execution. A delay statement delays the execution of a task for a specified duration or until a specified time. An entry call statement is written as a procedure call statement; it requests an operation on a task or on a protected object, blocking the caller until the operation can be performed. A called task may accept an entry call by executing a corresponding accept statement, which specifies the actions then to be performed as part of the rendezvous with the calling task. An entry call on a protected object is processed when the corresponding entry barrier evaluates to true, whereupon the body of the entry is executed. The requeue statement permits the provision of a service as a number of related activities with preference control. One form of the select statement allows a selective wait for one of several alternative rendezvous. Other forms of the select statement allow conditional or timed entry calls and the asynchronous transfer of control in response to some triggering event.

Execution of a program unit may encounter error situations in which normal program execution cannot continue. For example, an arithmetic computation may exceed the maximum allowed value of a number, or an attempt may be made to access an array component by using an incorrect index value. To deal with such error situations, the statements of a program unit can be textually followed by exception handlers that specify the actions to be taken when the error situation arises. Exceptions can be raised explicitly by a raise statement.

#### Data Types

Every object in the language has a type, which characterizes a set of values and a set of applicable operations. The main classes of types are elementary types (comprising enumeration, numeric, and access types) and composite types (including array and record types).

An enumeration type defines an ordered set of distinct enumeration literals, for example a list of states or an alphabet of characters. The enumeration types Boolean, Character, Wide\_Character, and Wide\_Wide\_Character are predefined.

Numeric types provide a means of performing exact or approximate numerical computations. Exact computations use integer types, which denote sets of consecutive integers. Approximate computations use either fixed point types, with absolute bounds on the error, or floating point types, with relative bounds on the error. The numeric types Integer, Float, and Duration are predefined.

Composite types allow definitions of structured objects with related components. The composite types in the language include arrays and records. An array is an object with indexed components of the same type. A record is an object with named components of possibly different types. Task and protected types are also forms of composite types. The array types String, Wide\_String, and Wide\_Wide\_String are predefined.

Record, task, and protected types may have special components called discriminants which parameterize the type. Variant record structures that depend on the values of discriminants can be defined within a record type.

Access types allow the construction of linked data structures. A value of an access type represents a reference to an object declared as aliased or to an object created by the evaluation of an allocator. Several variables of an access type may designate the same object, and components of one object may designate the same or other objects. Both the elements in such linked data structures and their relation to other elements can be altered during program execution. Access types also permit references to subprograms to be stored, passed as parameters, and ultimately dereferenced as part of an indirect call.

Private types permit restricted views of a type. A private type can be defined in a package so that only the logically necessary properties are made visible to the users of the type. The full structural details that are externally irrelevant are then only available within the package and any child units.

From any type a new type may be defined by derivation. A type, together with its derivatives (both direct and indirect) form a derivation class. Class-wide operations may be defined that accept as a parameter an operand of any type in a derivation class. For record and private types, the derivatives

may be extensions of the parent type. Types that support these object-oriented capabilities of class-wide operations and type extension must be tagged, so that the specific type of an operand within a derivation class can be identified at run time. When an operation of a tagged type is applied to an operand whose specific type is not known until run time, implicit dispatching is performed based on the tag of the operand.

Interface types provide abstract models from which other interfaces and types may be composed and derived. This provides a reliable form of multiple inheritance. Interface types may also be implemented by task types and protected types thereby enabling concurrent programming and inheritance to be merged.

The concept of a type is further refined by the concept of a subtype, whereby a user can constrain the set of allowed values of a type. Subtypes can be used to define subranges of scalar types, arrays with a limited set of index values, and records and private types with particular discriminant values.

#### Other Facilities

Aspect clauses can be used to specify the mapping between types and features of an underlying machine. For example, the user can specify that objects of a given type must be represented with a given number of bits, or that the components of a record are to be represented using a given storage layout. Other features allow the controlled use of low level, nonportable, or implementation-dependent aspects, including the direct insertion of machine code.

The predefined environment of the language provides for input-output and other capabilities by means of standard library packages. Input-output is supported for values of user-defined as well as of predefined types. Standard means of representing values in display form are also provided.

The predefined standard library packages provide facilities such as string manipulation, containers of various kinds (vectors, lists, maps, etc.), mathematical functions, random number generation, and access to the execution environment.

The specialized annexes define further predefined library packages and facilities with emphasis on areas such as real-time scheduling, interrupt handling, distributed systems, numerical computation, and high-integrity systems.

Finally, the language provides a powerful means of parameterization of program units, called generic program units. The generic parameters can be types and subprograms (as well as objects and packages) and so allow general algorithms and data structures to be defined that are applicable to all types of a given class.

#### **Language Changes**

This International Standard replaces the second edition of 1995. It modifies the previous edition by making changes and additions that improve the capability of the language and the reliability of programs written in the language. This edition incorporates the changes from Amendment 1 (ISO/IEC 8652:1995:AMD 1:2007), which were designed to improve the portability of programs, interfacing to other languages, and both the object-oriented and real-time capabilities.

Significant changes originating in Amendment 1 are incorporated:

- Support for program text is extended to cover the entire ISO/IEC 10646:2003 repertoire. Execution support now includes the 32-bit character set. See subclauses 2.1, 3.5.2, 3.6.3, A.1, A.3, and A.4.
- The object-oriented model has been improved by the addition of an interface facility which provides multiple inheritance and additional flexibility for type extensions. See subclauses 3.4, 3.9, and 7.3. An alternative notation for calling operations more akin to that used in other languages has also been added. See subclause 4.1.3.
- Access types have been further extended to unify properties such as the ability to access constants and to exclude null values. See clause 3.10. Anonymous access types are now

permitted more freely and anonymous access-to-subprogram types are introduced. See subclauses 3.3, 3.6, 3.10, and 8.5.1.

- The control of structure and visibility has been enhanced to permit mutually dependent references between units and finer control over access from the private part of a package. See subclauses 3.10.1 and 10.1.2. In addition, limited types have been made more useful by the provision of aggregates, constants, and constructor functions. See subclauses 4.3, 6.5, and 7.5.
- The predefined environment has been extended to include additional time and calendar operations, improved string handling, a comprehensive container library, file and directory management, and access to environment variables. See subclauses 9.6.1, A.4, A.16, A.17, and A.18.
- Two of the Specialized Needs Annexes have been considerably enhanced:
  - The Real-Time Systems Annex now includes the Ravenscar profile for high-integrity systems, further dispatching policies such as Round Robin and Earliest Deadline First, support for timing events, and support for control of CPU time utilization. See subclauses D.2, D.13, D.14, and D.15.
  - The Numerics Annex now includes support for real and complex vectors and matrices as previously defined in ISO/IEC 13813:1997 plus further basic operations for linear algebra. See subclause G.3.
- The overall reliability of the language has been enhanced by a number of improvements. These include new syntax which detects accidental overloading, as well as pragmas for making assertions and giving better control over the suppression of checks. See subclauses 6.1, 11.4.2, and 11.5.

In addition, this third edition makes enhancements to address two important issues, namely, the particular problems of multiprocessor architectures, and the need to further increase the capabilities regarding assertions for correctness. It also makes additional changes and additions that improve the capability of the language and the reliability of programs written in the language.

The following significant changes with respect to the 1995 edition as amended by Amendment 1 are incorporated:

- New syntax (the aspect specification) is introduced to enable properties to be specified for various entities in a more structured manner than through pragmas. See subclause 13.1.1.
- The concept of assertions introduced in the 2005 edition is extended with the ability to specify preconditions and postconditions for subprograms, and invariants for private types. The concept of constraints in defining subtypes is supplemented with subtype predicates that enable subsets to be specified other than as simple ranges. These properties are all indicated using aspect specifications. See subclauses 3.2.4, 6.1.1, and 7.3.2.
- New forms of expressions are introduced. These are if expressions, case expressions, quantified expressions, and expression functions. As well as being useful for programming in general by avoiding the introduction of unnecessary assignments, they are especially valuable in conditions and invariants since they avoid the need to introduce auxiliary functions. See subclauses 4.5.7, 4.5.8, and 6.8. Membership tests are also made more flexible. See subclauses 4.4 and 4.5.2.
- A number of changes are made to subprogram parameters. Functions may now have parameters of all modes. In order to mitigate consequent (and indeed existing) problems of inadvertent order dependence, rules are introduced to reduce aliasing. A parameter may now be explicitly marked as aliased and the type of a parameter may be incomplete in certain circumstances. See subclauses 3.10.1, 6.1, and 6.4.1.
- The use of access types is now more flexible. The rules for accessibility and certain conversions are improved. See subclauses 3.10.2, 4.5.2, 4.6, and 8.6. Furthermore, better control of storage pools is provided. See subclause 13.11.4.

- The Real-Time Systems Annex now includes facilities for defining domains of processors and assigning tasks to them. Improvements are made to scheduling and budgeting facilities. See subclauses D.10.1, D.14, and D.16.
- A number of important improvements are made to the standard library. These include packages for conversions between strings and UTF encodings, and classification functions for wide and wide characters. Internationalization is catered for by a package giving locale information. See subclauses A.3, A.4.11, and A.19. The container library is extended to include bounded forms of the existing containers and new containers for indefinite objects, multiway trees, and queues. See subclause A.18.
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  i iteration o id 5.5.2. • Finally, certain features are added primarily to ease the use of containers, such as the ability to iterate over all elements in a container without having to encode the iteration. These can also be used for iteration over arrays, and within quantified expressions. See subclauses 4.1.5, 4.1.6, 5.5.1, and 5.5.2.

#### Instructions for Comment Submission

Informal comments on this International Standard may be sent via e-mail to **ada-comment@ada-auth.org**. If appropriate, the Project Editor will initiate the defect correction procedure.

Comments should use the following format:

!topic Title summarizing comment !reference Ada 2012 RMss.ss(pp) !from Author Name yy-mm-dd !keywords keywords related to topic !discussion

text of discussion

where ss.ss is the clause or subclause number, pp is the paragraph number where applicable, and yymm-dd is the date the comment was sent. The date is optional, as is the !keywords line.

Please use a descriptive "Subject" in your e-mail message, and limit each message to a single comment.

When correcting typographical errors or making minor wording suggestions, please put the correction directly as the topic of the comment; use square brackets [] to indicate text to be omitted and curly braces {} to indicate text to be added, and provide enough context to make the nature of the suggestion self-evident or put additional information in the body of the comment, for example:

!topic [c]{C}haracter
!topic it[']s meaning is not defined

Formal requests for interpretations and for reporting defects in this International Standard may be made in accordance with the ISO/IEC JTC 1 Directives and the ISO/IEC JTC 1/SC 22 policy for interpretations. National Bodies may submit a Defect Report to ISO/IEC JTC 1/SC 22 for resolution under the JTC 1 procedures. A response will be provided and, if appropriate, a Technical Corrigendum will be issued in accordance with the procedures.

# Information technology — Programming Languages — Ada

#### 1 General

### 1.1 Scope

This International Standard specifies the form and meaning of programs written in Ada. Its purpose is to promote the portability of Ada programs to a variety of computing systems.

Ada is a programming language designed to support the construction of long-lived, highly reliable software systems. The language includes facilities to define packages of related types, objects, and operations. The packages may be parameterized and the types may be extended to support the construction of libraries of reusable, adaptable software components. The operations may be implemented as subprograms using conventional sequential control structures, or as entries that include synchronization of concurrent threads of control as part of their invocation. Ada supports object-oriented programming by providing classes and interfaces, inheritance, polymorphism of variables and methods, and generic units. The language treats modularity in the physical sense as well, with a facility to support separate compilation.

The language provides rich support for real-time, concurrent programming, and includes facilities for multicore and multiprocessor programming. Errors can be signaled as exceptions and handled explicitly. The language also covers systems programming; this requires precise control over the representation of data and access to system-dependent properties. Finally, a predefined environment of standard packages is provided, including facilities for, among others, input-output, string manipulation, numeric elementary functions, and random number generation, and definition and use of containers.

#### **1.1.1 Extent**

This International Standard specifies:

- The form of a program written in Ada;
- The effect of translating and executing such a program;
- The manner in which program units may be combined to form Ada programs;
- The language-defined library units that a conforming implementation is required to supply;
- The permissible variations within the standard, and the manner in which they are to be documented:

- Those violations of the standard that a conforming implementation is required to detect, and the effect of attempting to translate or execute a program containing such violations;
- Those violations of the standard that a conforming implementation is not required to detect.

This International Standard does not specify:

- The means whereby a program written in Ada is transformed into object code executable by a processor:
- The means whereby translation or execution of programs is invoked and the executing units are controlled;
- The size or speed of the object code, or the relative execution speed of different language constructs:
- The form or contents of any listings produced by implementations; in particular, the form or contents of error or warning messages;
- The effect of unspecified execution.
- The size of a program or program unit that will exceed the capacity of a particular conforming implementation.

#### 1.1.2 Structure

This International Standard contains thirteen clauses, fifteen annexes, and an index.

The core of the Ada language consists of:

- Clauses 1 through 13
- Annex A, "Predefined Language Environment"
- Annex B, "Interface to Other Languages"
- Annex J, "Obsolescent Features"

The following Specialized Needs Annexes define features that are needed by certain application areas:

- Annex C, "Systems Programming"
- Annex D, "Real-Time Systems"
- Annex E, "Distributed Systems"
- Annex F, "Information Systems"
- Annex G, "Numerics"
- Annex H, "High Integrity Systems"

The core language and the Specialized Needs Annexes are normative, except that the material in each of the items listed below is informative:

- Text under a NOTES or Examples heading.
- Each subclause whose title starts with the word "Example" or "Examples".

All implementations shall conform to the core language. In addition, an implementation may conform separately to one or more Specialized Needs Annexes.

The following Annexes are informative:

- Annex K, "Language-Defined Aspects and Attributes"
- Annex L, "Language-Defined Pragmas"
- Annex M, "Summary of Documentation Requirements"
- Annex N, "Glossary"

- Annex P, "Syntax Summary"
- Annex Q, "Language-Defined Entities"

Each section is divided into subclauses that have a common structure. Each clause and subclause first introduces its subject. After the introductory text, text is labeled with the following headings:

Syntax

Syntax rules (indented).

Name Resolution Rules

Compile-time rules that are used in name resolution, including overload resolution.

Legality Rules

Rules that are enforced at compile time. A construct is *legal* if it obeys all of the Legality Rules.

Static Semantics

A definition of the compile-time effect of each construct.

Post-Compilation Rules

Rules that are enforced before running a partition. A partition is legal if its compilation units are legal and it obeys all of the Post-Compilation Rules.

**Dynamic Semantics** 

A definition of the run-time effect of each construct.

Bounded (Run-Time) Errors

Situations that result in bounded (run-time) errors (see 1.1.5).

Erroneous Execution

Situations that result in erroneous execution (see 1.1.5).

Implementation Requirements

Additional requirements for conforming implementations.

Documentation Requirements

Documentation requirements for conforming implementations.

Metrics

Metrics that are specified for the time/space properties of the execution of certain language constructs.

Implementation Permissions

Additional permissions given to the implementer.

Implementation Advice

Optional advice given to the implementer. The word "should" is used to indicate that the advice is a recommendation, not a requirement. It is implementation defined whether or not a given recommendation is obeyed.

NOTES

1 Notes emphasize consequences of the rules described in the (sub)clause or elsewhere. This material is informative.

Examples

Examples illustrate the possible forms of the constructs described. This material is informative.

### 1.1.3 Conformity of an Implementation with the Standard

Implementation Requirements

A conforming implementation shall:

- Translate and correctly execute legal programs written in Ada, provided that they are not so large as to exceed the capacity of the implementation;
- Identify all programs or program units that are so large as to exceed the capacity of the implementation (or raise an appropriate exception at run time);
- Identify all programs or program units that contain errors whose detection is required by this International Standard;
- Supply all language-defined library units required by this International Standard;
- Contain no variations except those explicitly permitted by this International Standard, or those that are impossible or impractical to avoid given the implementation's execution environment;
- Specify all such variations in the manner prescribed by this International Standard.

The *external effect* of the execution of an Ada program is defined in terms of its interactions with its external environment. The following are defined as *external interactions*:

- Any interaction with an external file (see A.7);
- The execution of certain code\_statements (see 13.8); which code\_statements cause external interactions is implementation defined.
- Any call on an imported subprogram (see Annex B), including any parameters passed to it;
- Any result returned or exception propagated from a main subprogram (see 10.2) or an exported subprogram (see Annex B) to an external caller;
- Any read or update of an atomic or volatile object (see C.6);
- The values of imported and exported objects (see Annex B) at the time of any other interaction with the external environment.

A conforming implementation of this International Standard shall produce for the execution of a given Ada program a set of interactions with the external environment whose order and timing are consistent with the definitions and requirements of this International Standard for the semantics of the given program.

An implementation that conforms to this Standard shall support each capability required by the core language as specified. In addition, an implementation that conforms to this Standard may conform to one or more Specialized Needs Annexes (or to none). Conformance to a Specialized Needs Annex means that each capability required by the Annex is provided as specified.

An implementation conforming to this International Standard may provide additional aspects, attributes, library units, and pragmas. However, it shall not provide any aspect, attribute, library unit, or pragma having the same name as an aspect, attribute, library unit, or pragma (respectively) specified in a Specialized Needs Annex unless the provided construct is either as specified in the Specialized Needs Annex or is more limited in capability than that required by the Annex. A program that attempts to use an unsupported capability of an Annex shall either be identified by the implementation before run time or shall raise an exception at run time.

Documentation Requirements

Certain aspects of the semantics are defined to be either *implementation defined* or *unspecified*. In such cases, the set of possible effects is specified, and the implementation may choose any effect in the set. Implementations shall document their behavior in implementation-defined situations, but

documentation is not required for unspecified situations. The implementation-defined characteristics are summarized in M.2.

The implementation may choose to document implementation-defined behavior either by documenting what happens in general, or by providing some mechanism for the user to determine what happens in a particular case.

Implementation Advice

If an implementation detects the use of an unsupported Specialized Needs Annex feature at run time, it should raise Program Error if feasible.

If an implementation wishes to provide implementation-defined extensions to the functionality of a language-defined library unit, it should normally do so by adding children to the library unit.

NOTES

2 The above requirements imply that an implementation conforming to this Standard may support some of the capabilities required by a Specialized Needs Annex without supporting all required capabilities.

### 1.1.4 Method of Description and Syntax Notation

The form of an Ada program is described by means of a context-free syntax together with context-dependent requirements expressed by narrative rules.

The meaning of Ada programs is described by means of narrative rules defining both the effects of each construct and the composition rules for constructs.

The context-free syntax of the language is described using a simple variant of Backus-Naur Form. In particular:

• Lower case words in a sans-serif font, some containing embedded underlines, are used to denote syntactic categories, for example:

```
case_statement
```

• Boldface words are used to denote reserved words, for example:

```
array
```

• Square brackets enclose optional items. Thus the two following rules are equivalent.

```
simple_return_statement ::= return [expression];
simple_return_statement ::= return; | return expression;
```

• Curly brackets enclose a repeated item. The item may appear zero or more times; the repetitions occur from left to right as with an equivalent left-recursive rule. Thus the two following rules are equivalent.

```
term ::= factor {multiplying_operator factor}
term ::= factor | term multiplying_operator factor
```

• A vertical line separates alternative items unless it occurs immediately after an opening curly bracket, in which case it stands for itself:

```
constraint ::= scalar_constraint | composite_constraint
discrete_choice | discrete_choice { | discrete_choice }
```

• If the name of any syntactic category starts with an italicized part, it is equivalent to the category name without the italicized part. The italicized part is intended to convey some semantic information. For example <code>subtype\_name</code> and <code>task\_name</code> are both equivalent to name alone.

The delimiters, compound delimiters, reserved words, and numeric\_literals are exclusively made of the characters whose code point is between 16#20# and 16#7E#, inclusively. The special characters for which names are defined in this International Standard (see 2.1) belong to the same range. For